

# PART 1

-

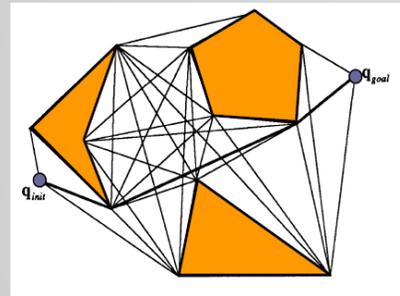
# MOTION PLANNING

MOTION PLANNING ALGORITHMS



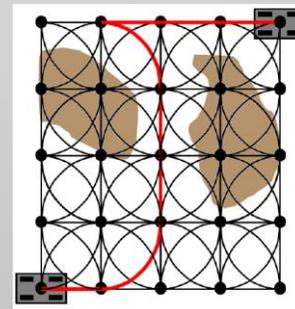
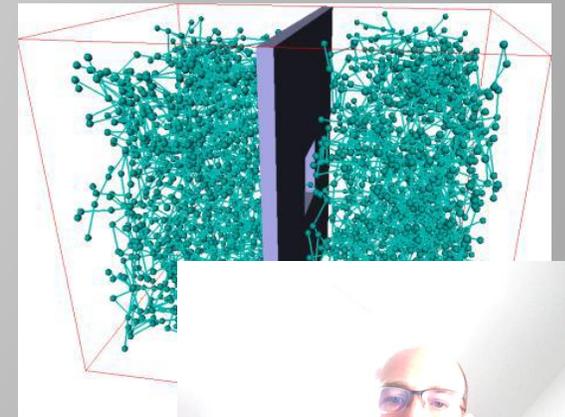
# OVERVIEW

- Roadmap methods
  - Visibility graphs
  - Voronoi diagrams
- Decomposition methods
  - Approximate
  - Exact
- Potential fields
- Sampling-based Planning
  - PRM
  - RRT



- Deterministic, heuristic or probabilistic
- Planning algorithms evaluation criteria:

- Completeness
- Optimality
- Speed
- Generality

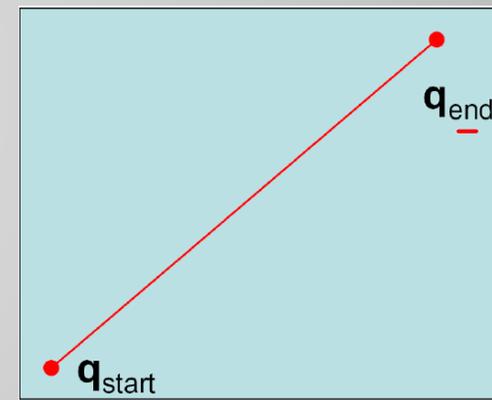


- Applications to Computer Animation
- Dynamic environments
- Crowd Simulation



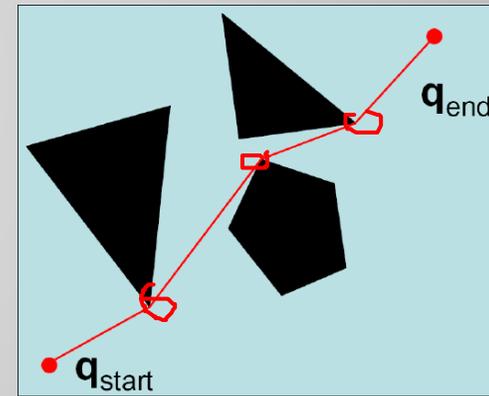
# VISIBILITY GRAPHS

IN THE ABSENCE OF OBSTACLES, THE BEST  
PATH IS THE STRAIGHT LINE BETWEEN  
 $Q_{\text{START}}$  AND  $Q_{\text{END}}$



# VISIBILITY GRAPHS

ASSUMING POLYGONAL OBSTACLES:  
IT LOOKS LIKE THE SHORTEST PATH IS A  
SEQUENCE OF STRAIGHT LINES JOINING  
THE VERTICES OF THE OBSTACLES



# VISIBILITY GRAPHS

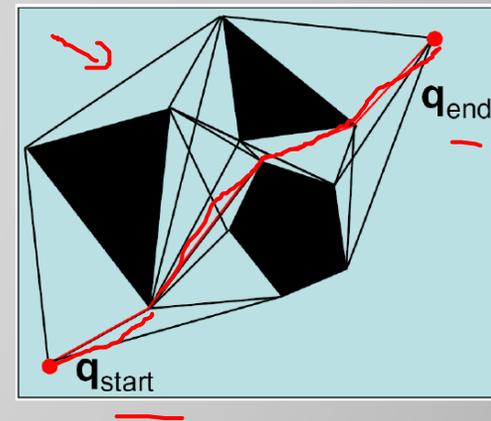
VISIBILITY GRAPH  $G$

= SET OF UNBLOCKED LINES BETWEEN  
THE VERTICES OF THE OBSTACLES

+  $Q_{\text{START}}$  AND  $Q_{\text{GOAL}}$

A NODE  $P$  IS LINKED TO A NODE  $P'$  IF  $P'$   
VISIBLE FROM  $P$

SOLUTION = SHORTEST PATH IN THE  
VISIBILITY GRAPH



# VISIBILITY GRAPHS

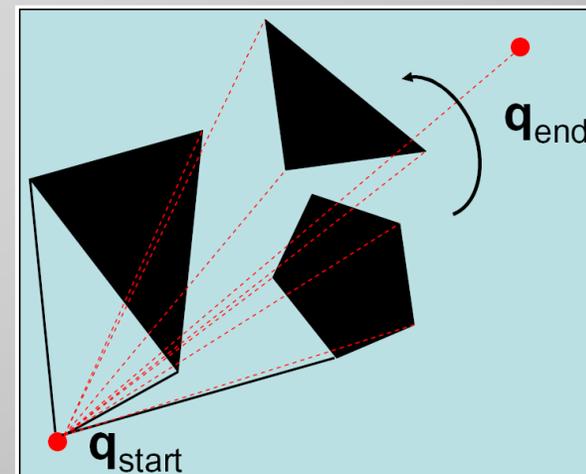
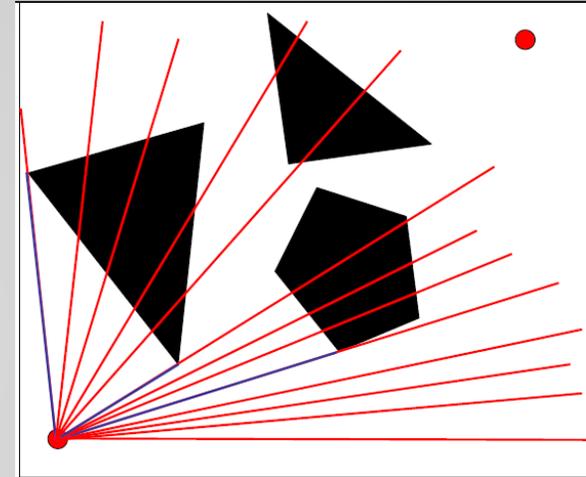
CONSTRUCTION: SWEEP ALGORITHM

SWEEP A LINE ORIGINATING AT EACH VERTEX

RECORD THOSE LINES THAT END AT VISIBLE VERTICES

COMPLEXITY

- $N$  = total number of vertices of the obstacle polygons
- Naïve:  $O(N^3)$
- Sweep:  $O(N^2 \log N)$

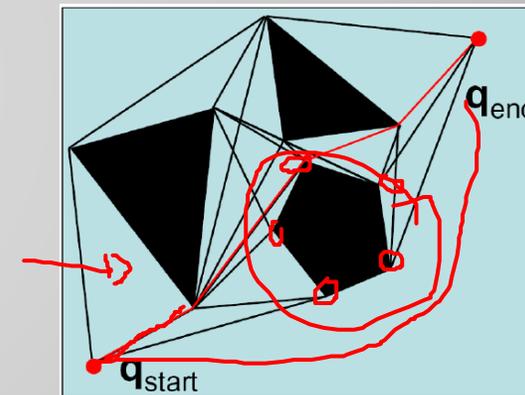


# VISIBILITY GRAPHS

## SHORTEST PATH BUT:

- Tries to stay as close as possible to obstacles
- Any execution error will lead to a collision
- Complicated in  $\gg 2$  dimensions

WE MAY NOT CARE ABOUT STRICT OPTIMALITY  
SO LONG AS WE FIND A SAFE PATH. STAYING  
AWAY FROM OBSTACLES IS MORE IMPORTANT  
THAN FINDING THE SHORTEST PATH



→ NEED TO DEFINE OTHER TYPES OF “ROADMAPS”

Lozano-Pérez, Tomás; Wesley, Michael A. (1979), "An algorithm for planning collision-free paths among obstacles", *Communications of the ACM*, **22** (10): 560–570, [doi:10.1145/359156.359164](https://doi.org/10.1145/359156.359164)

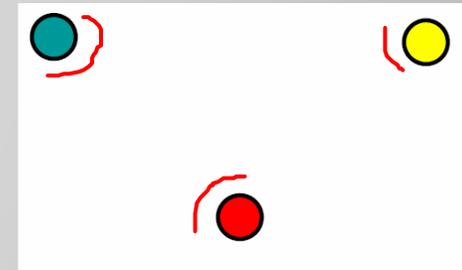


# VORONOI DIAGRAM

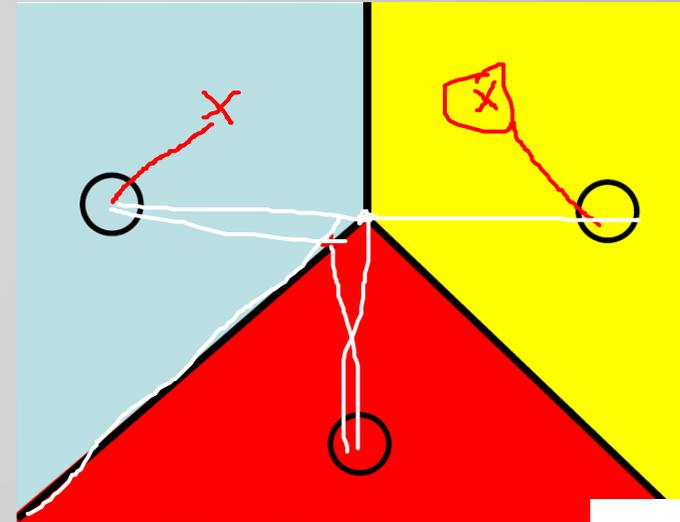
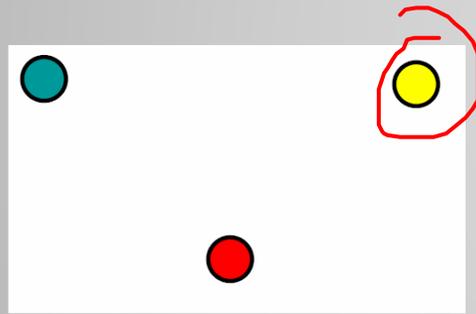
---

GIVEN A SET OF DATA POINTS IN THE PLANE:

- Color the entire plane such that the color of any point in the plane is the same as the color of its nearest



# VORONOI DIAGRAM



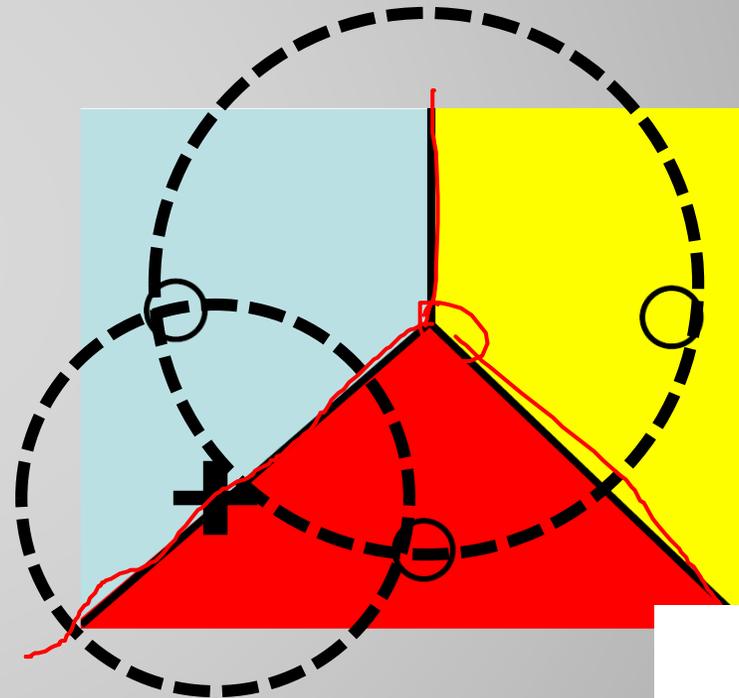
# VORONOI DIAGRAM

VORONOI DIAGRAM:

THE SET OF LINE SEGMENTS SEPARATING  
THE REGIONS CORRESPONDING TO  
DIFFERENT COLORS

LINE SEGMENT = POINTS EQUIDISTANT  
FROM 2 DATA POINTS

VERTICES = POINTS EQUIDISTANT FROM >  
2 DATA POINTS



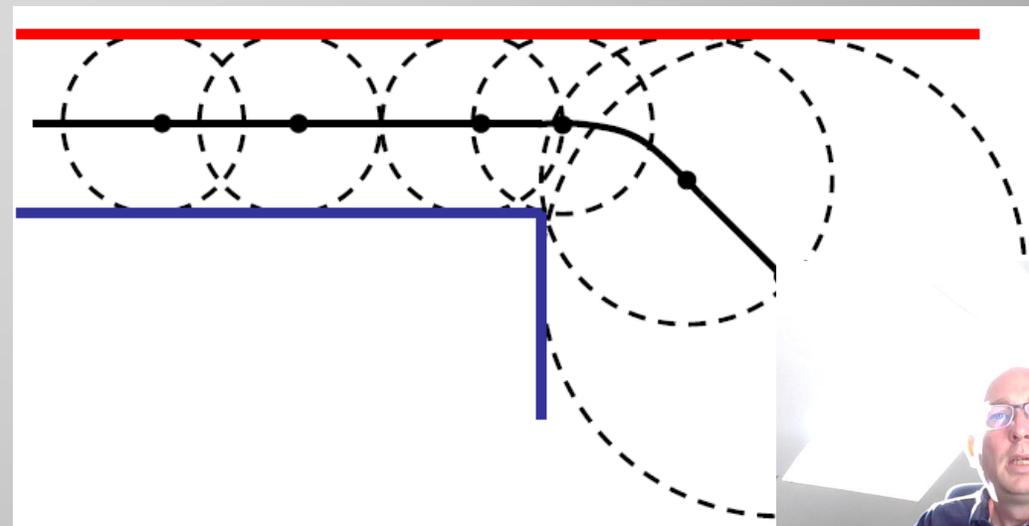
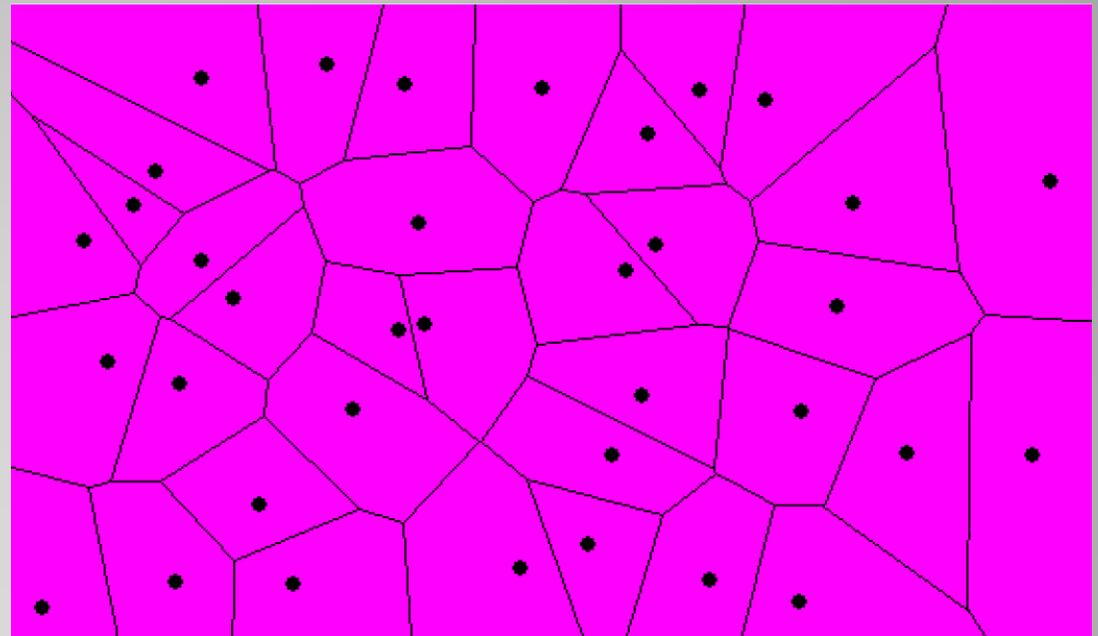
# VORONOI

## COMPLEXITY (IN THE PLANE):

- $O(N \log N)$  time
- $O(N)$  space

## BEYOND POINTS:

- Edges are combinations of straight line segments and segments of quadratic curves
- Straight edges: Points equidistant from 2 lines
- Curved edges: Points equidistant from one corner and one line



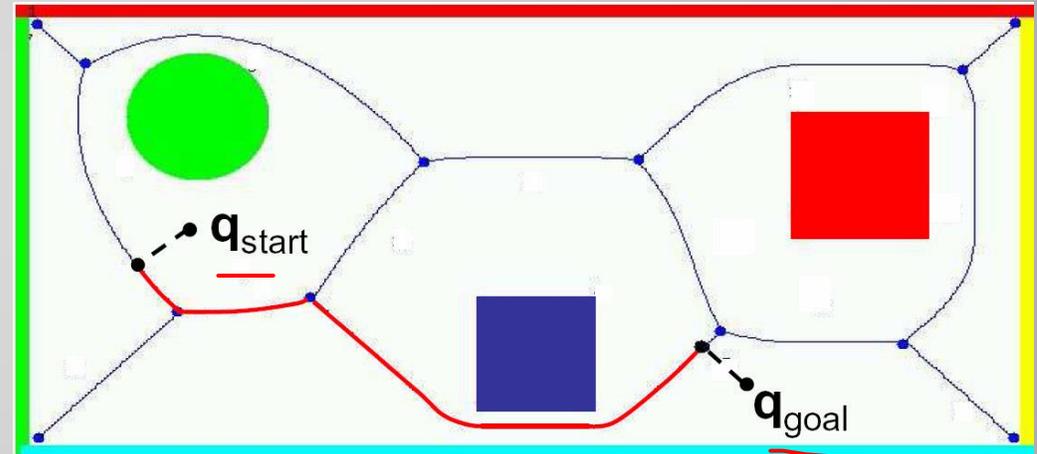
# VORONOI DIAGRAM

## KEY PROPERTY:

THE POINTS ON THE EDGES OF THE VORONOI DIAGRAM ARE THE *FURTHEST* FROM THE OBSTACLES

## IDEA:

CONSTRUCT A PATH BETWEEN  $q_{\text{START}}$  AND  $q_{\text{GOAL}}$  BY FOLLOWING EDGES ON THE VORONOI DIAGRAM (USE THE VORONOI DIAGRAM AS A ROADMAP GRAPH INSTEAD OF THE VISIBILITY GRAPH)



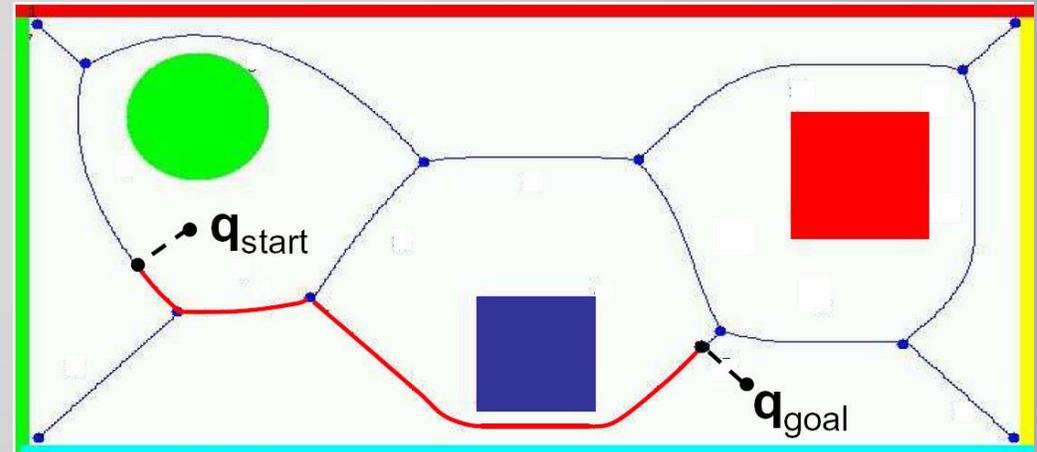
# VORONOI DIAGRAM

DIFFICULT TO COMPUTE IN HIGHER DIMENSIONS  
OR NONPOLYGONAL WORLDS

- Approximate algorithms exist
- Use of Voronoi is not necessarily the best

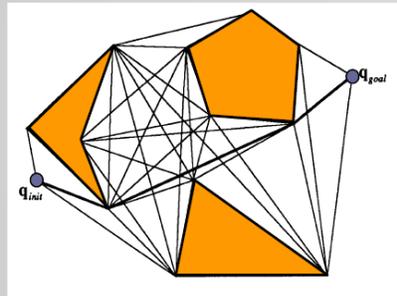
(HEURISTIC (“STAY AWAY FROM OBSTACLES”))  
CAN LEAD TO PATHS THAT ARE MUCH TOO  
CONSERVATIVE

CAN BE UNSTABLE: SMALL CHANGES IN  
OBSTACLE CONFIGURATION CAN LEAD TO LARGE  
CHANGES IN THE DIAGRAM

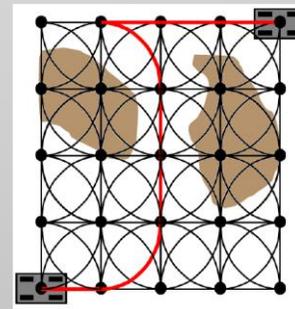
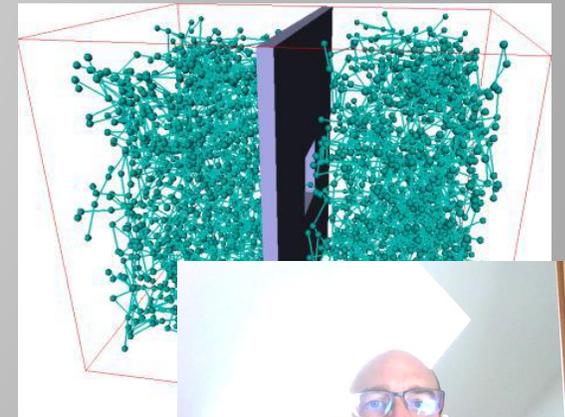


# OVERVIEW

- Roadmap methods
  - Visibility graphs
  - Voronoi diagrams
- Decomposition methods
  - Approximate
  - Exact
- Potential fields
- Sampling-based Planning
  - PRM
  - RRT
- Applications to Computer Animation
- Dynamic environments
- Crowd Simulation



- Deterministic, heuristic or probabilistic
- Planning algorithms evaluation criteria:
  - Completeness
  - Optimality
  - Speed
  - Generality



# APPROXIMATE CELL DECOMPOSITION

DEFINE A DISCRETE GRID IN  $C_{SPACE}$

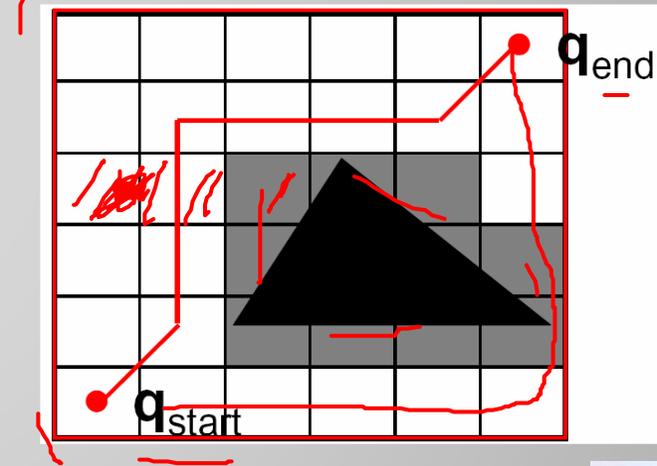
- Mark any cell of the grid that intersects  $C_{obs}$  as blocked

FIND PATH THROUGH REMAINING CELLS BY USING (FOR EXAMPLE) A\* (E.G., USE EUCLIDEAN DISTANCE AS HEURISTIC)

→ CANNOT BE COMPLETE AS DESCRIBED SO FAR. WHY?

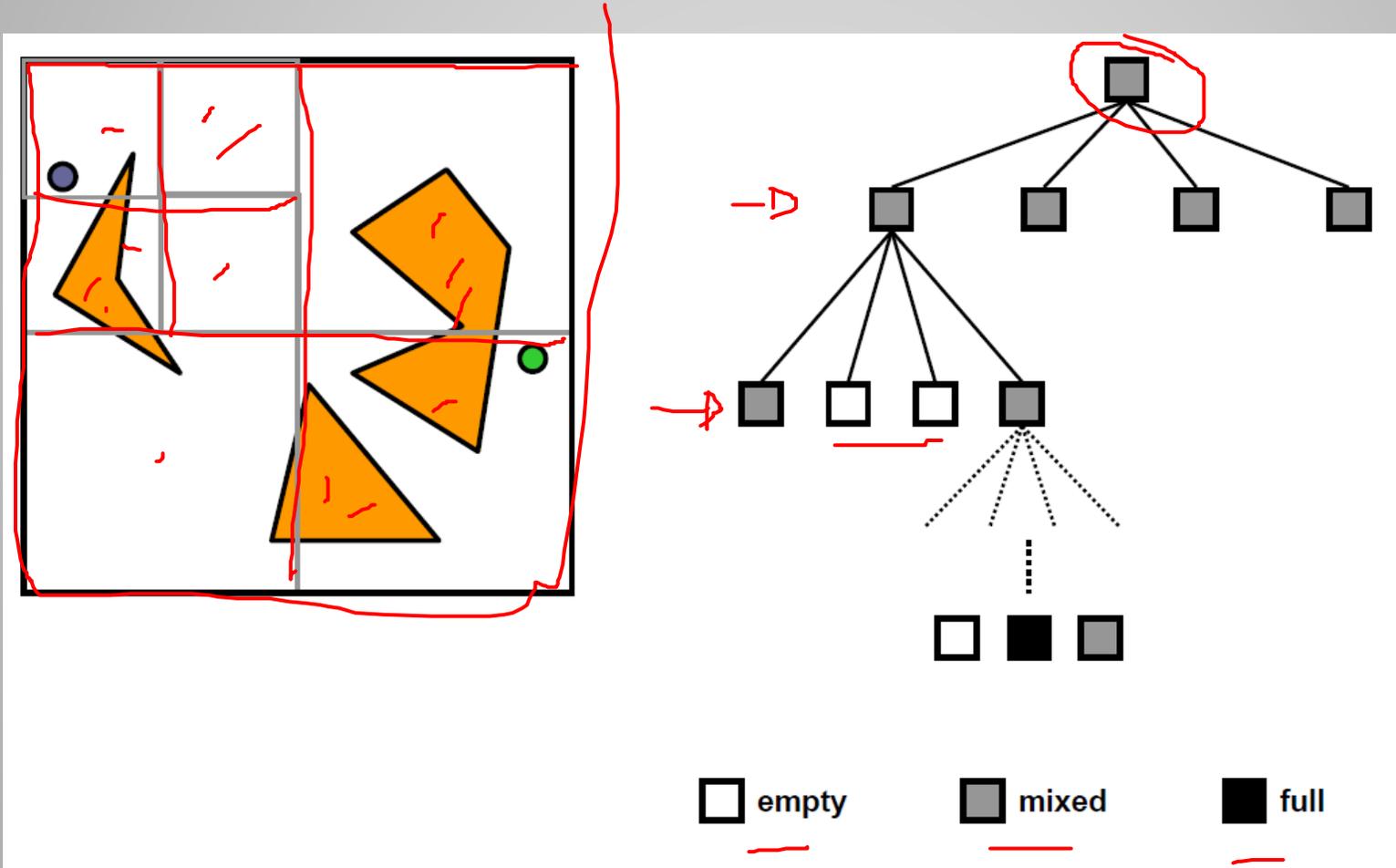
NOTE:

- deterministic approaches with exact decomposition are complete
- deterministic approaches with approximate decomposition are representation-complete

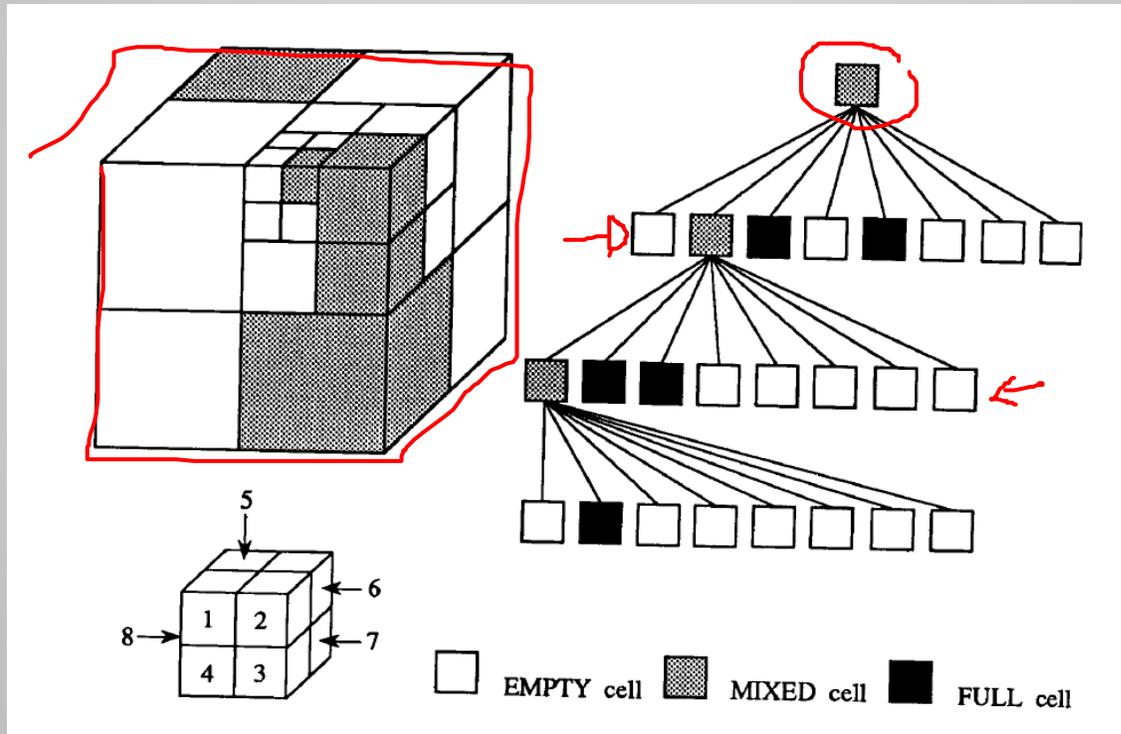




# QUADTREE DECOMPOSITION



# OCTREE DECOMPOSITION



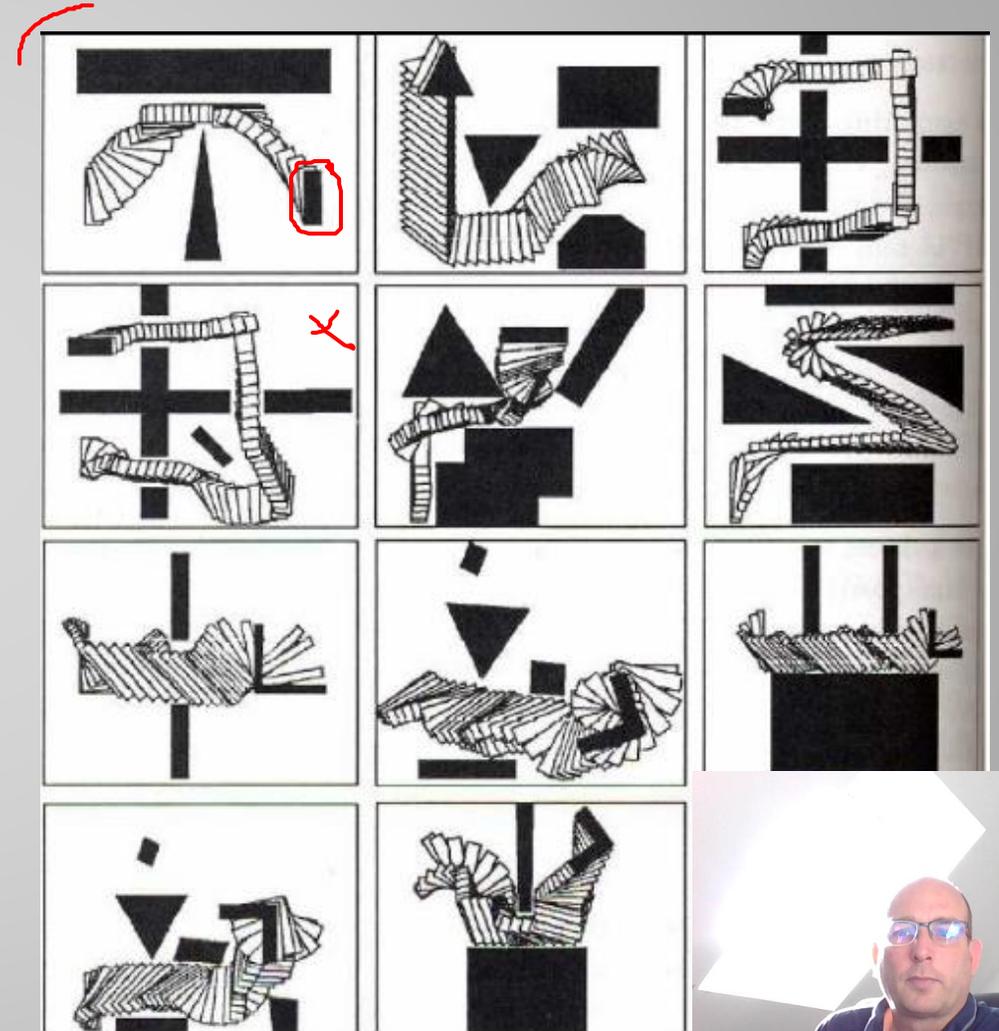
# APPROXIMATE CELL DECOMPOSITION

GOOD:

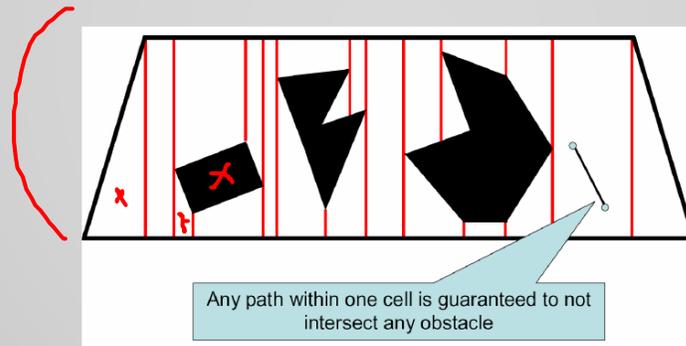
- Limited assumptions on obstacle configuration
- Approach used in practice
- Find obvious solutions quickly

BAD:

- No clear notion of optimality (“best” path)
- Trade-off completeness/computation
- Still difficult to use in high dimensions

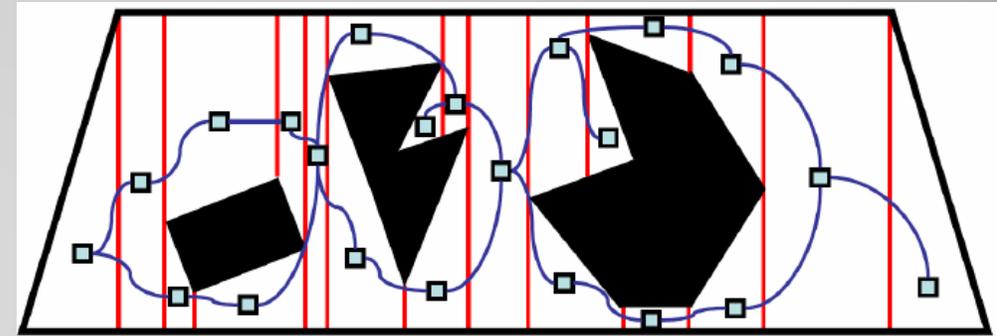


# EXACT CELL DECOMPOSITION

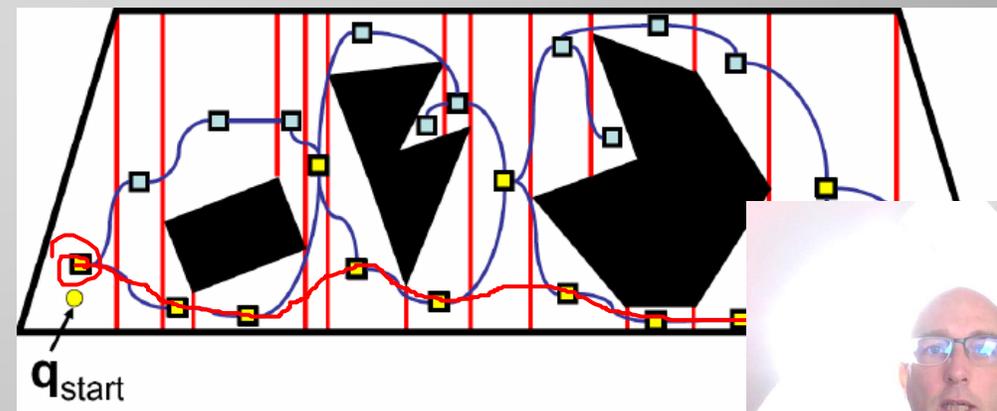


# EXACT CELL DECOMPOSITION

THE GRAPH OF CELLS DEFINES A  
ROADMAP



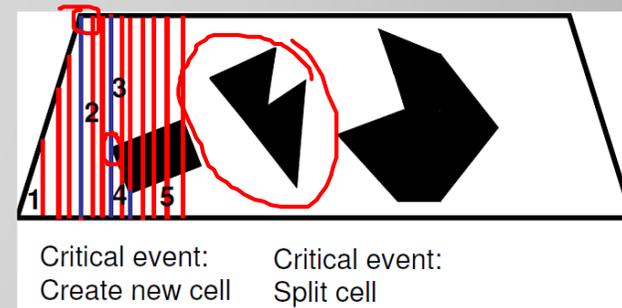
AND CAN BE USED TO FIND A PATH  
BETWEEN ANY TWO CONFIGURATION



# EXACT CELL DECOMPOSITION

## ( Plane Sweep algorithm

- Initialize current list of cells to empty
- Order the vertices of Cobs along the x direction
- For every vertex:
  - Construct the plane at the corresponding x location
  - Depending on the type of event:
- Split a current cell into 2 new cells OR
- Merge two of the current cells
  - Create a new cell
- Complexity (in 2-D):
  - Time:  $O(N \log N)$
  - Space:  $O(N)$



# OVERVIEW

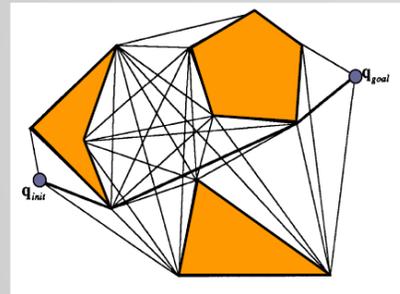
- Roadmap methods
  - Visibility graphs
  - Voronoi diagrams

## → Decomposition methods

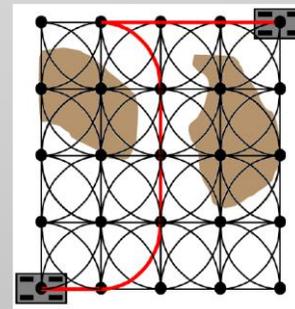
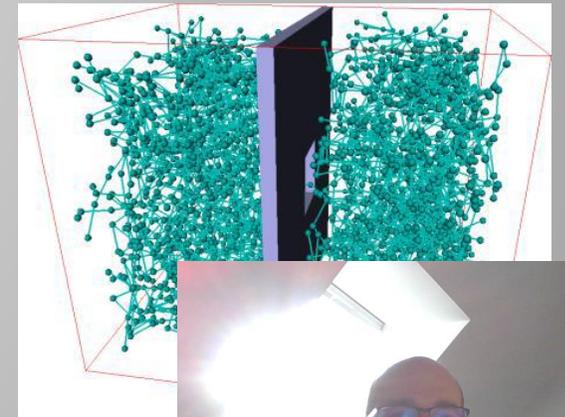
- Approximate
- Exact

## → Potential fields

- Sampling-based Planning
  - PRM
  - RRT
- Applications to Computer Animation
- Dynamic environments
- Crowd Simulation



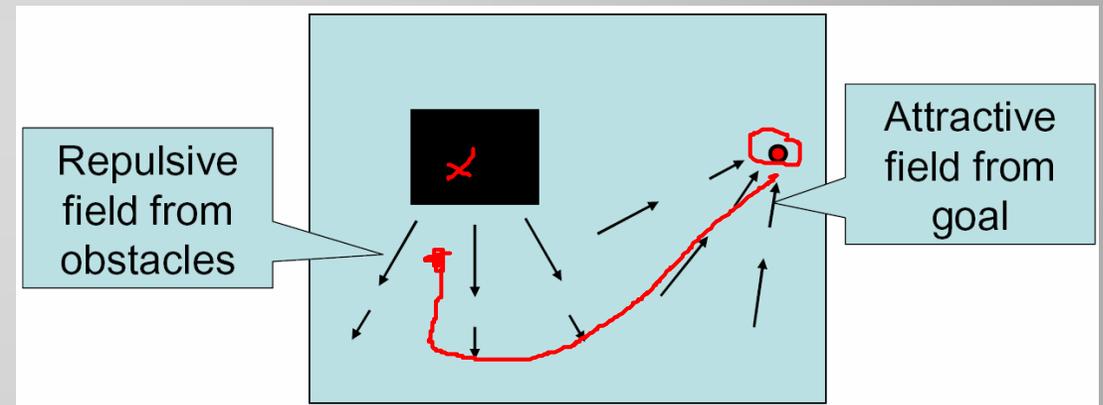
- Deterministic, heuristic or probabilistic
- Planning algorithms evaluation criteria:
  - • Completeness
  - Optimality
  - Speed
  - • Generality



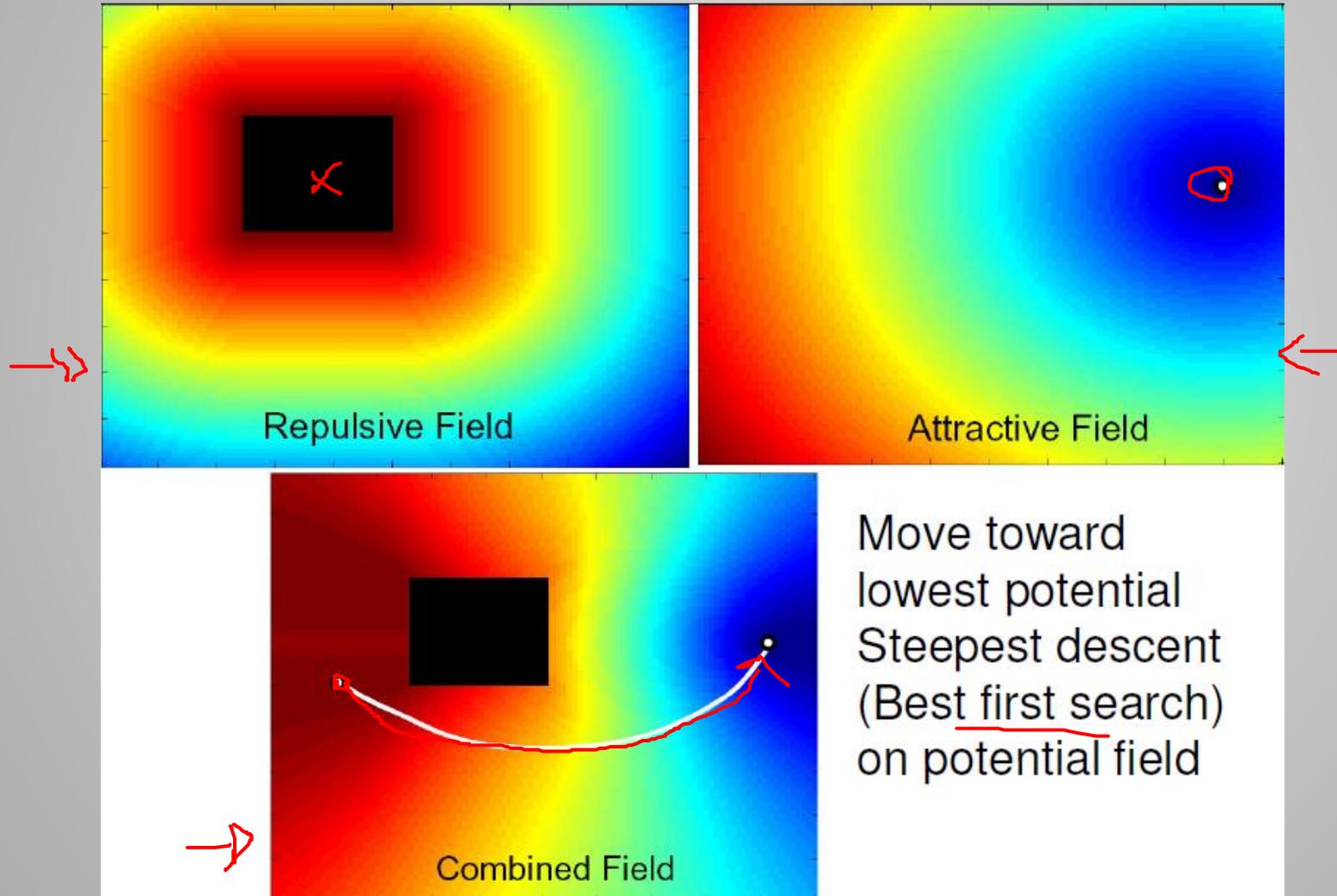
# POTENTIAL FIELDS

STAY AWAY FROM OBSTACLES: IMAGINE THAT THE OBSTACLES ARE MADE OF A MATERIAL THAT GENERATE A *REPULSIVE* FIELD

MOVE CLOSER TO THE GOAL: IMAGINE THAT THE GOAL LOCATION IS A PARTICLE THAT GENERATES AN ATTRACTIVE FIELD



# POTENTIAL FIELDS



# POTENTIAL FIELDS

$$\underline{U_g(\mathbf{q})} = d^2(\mathbf{q}, \underline{\mathbf{q}_{goal}})$$

Distance to goal state

$$\underline{U_o(\mathbf{q})} = \frac{1}{d^2(\mathbf{q}, \underline{Obstacles})}$$

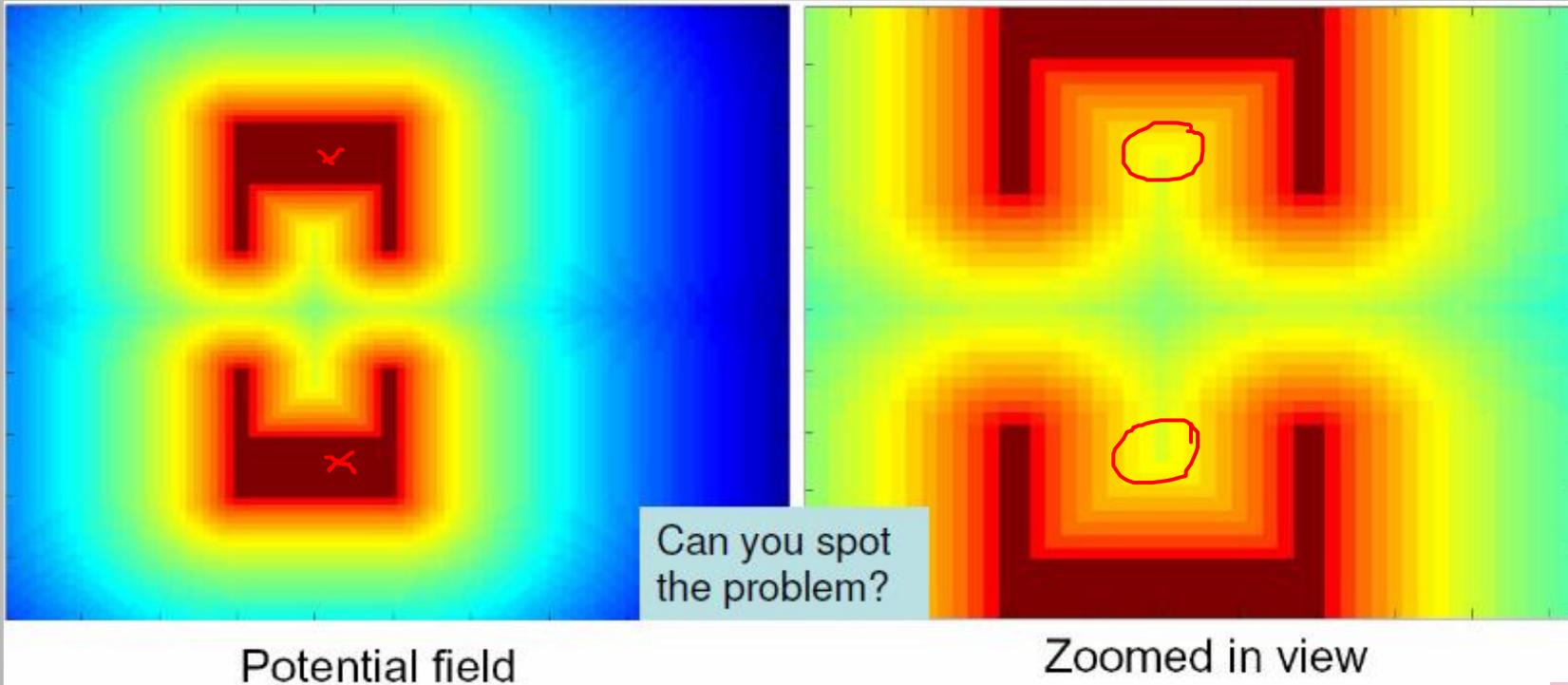
Distance to nearest obstacle point.  
Note: Can be computed efficiently by  
using the *distance transform*

$$\left( U(\mathbf{q}) = \underline{U_g(\mathbf{q})} + \lambda \underline{U_o(\mathbf{q})} \right)$$

$\lambda$  controls how far we  
stay from the obstacles



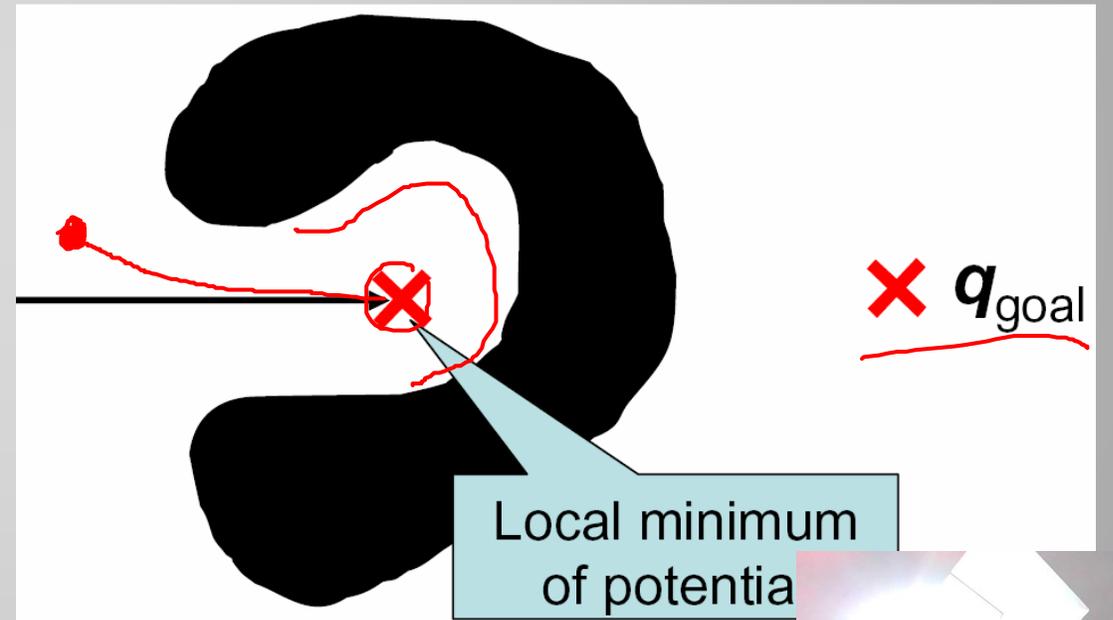
# POTENTIAL FIELDS



## POTENTIAL FIELDS IN GENERAL EXHIBIT LOCAL MINIMA

### SPECIAL CASE: NAVIGATION FUNCTION

- $U(\mathbf{q}_{\text{goal}}) = 0$
- For any  $\mathbf{q}$  different from  $\mathbf{q}_{\text{goal}}$ , there exists a neighbor  $\mathbf{q}'$  such that  $U(\mathbf{q}') < U(\mathbf{q})$



# GETTING OUT OF LOCAL MINIMA

Repeat

- If  $U(\mathbf{q}) = 0$  return Success
- If too many iterations return Failure
- Else:
  - Find neighbor  $\mathbf{q}_n$  of  $\mathbf{q}$  with smallest  $U(\mathbf{q}_n)$
  - If  $U(\mathbf{q}_n) < U(\mathbf{q})$  OR  $\mathbf{q}_n$  has not yet been visited
- Move to  $\mathbf{q}_n$  ( $\mathbf{q} \leftarrow \mathbf{q}_n$ )
- Remember  $\mathbf{q}_n$

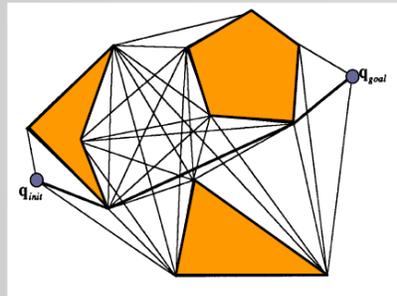
Repeat

- If  $U(\mathbf{q}) = 0$  return Success
- If too many iterations return Failure
- Else:
  - Find neighbor  $\mathbf{q}_n$  of  $\mathbf{q}$  with smallest  $U(\mathbf{q}_n)$
  - If  $U(\mathbf{q}_n) < U(\mathbf{q})$ 
    - Move to  $\mathbf{q}_n$  ( $\mathbf{q} \leftarrow \mathbf{q}_n$ )
  - Else
    - Take a random walk for  $T$  steps starting at  $\mathbf{q}_n$
    - Set  $\mathbf{q}$  to the configuration reached at the end of the random walk

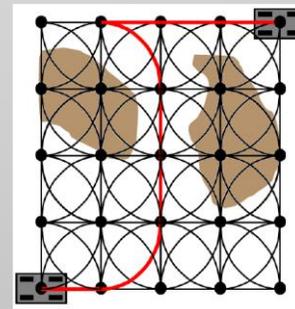
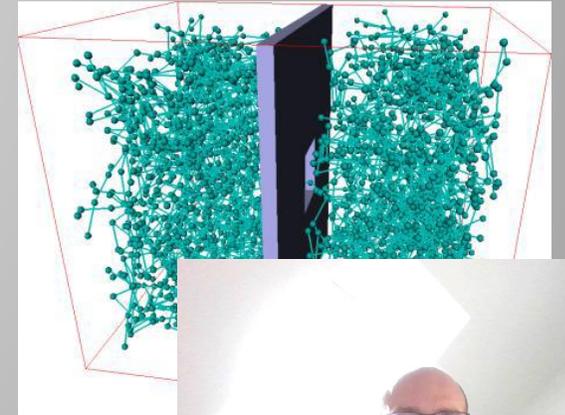


# OVERVIEW

- Roadmap methods
  - Visibility graphs
  - Voronoi diagrams
- Decomposition methods
  - Approximate
  - Exact
- Potential fields
- Sampling-based Planning
  - PRM
  - RRT
- Applications to Computer Animation
  - Dynamic environments
  - Crowd Simulation



- Deterministic, heuristic or probabilistic
- Planning algorithms evaluation criteria:
  - Completeness
  - Optimality
  - Speed
  - Generality



# PLANNING IN HIGH-DIMENSIONAL SPACES

## IDEAL: A COMPLETE PLANNER

- Guarantees to find a solution in finite time if exists
- Indicates the non existence of a solution in a finite time
- Note: completeness means the computation of 1 path, among others.

## PROBLEM: P-SPACE COMPLEXITY [REIF '79]

### SOLUTIONS:

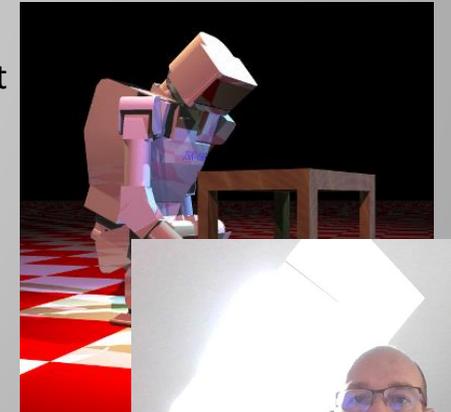
- Lower the dimension of the search space
- Limit the number of possible solutions (e.g., bound the search space)
- Sacrifice optimality
- Sacrifice completeness

KEY IDEA: INSTEAD OF SEARCHING THE WHOLE CONFIGURATION SPACE, RANDOMLY EXPLORE SOLUTIONS AND CAPTURE THEM

FACILITATES A « PROBING » INSTEAD OF EXHAUSTIVE EXPLORATION

### DRAWBACKS?

- Completeness and optimality are lost
- Classical trade-off however



# PROBABILISTIC ROADMAPS METHOD (PRM)

RELIES ON 3 ELEMENTS:

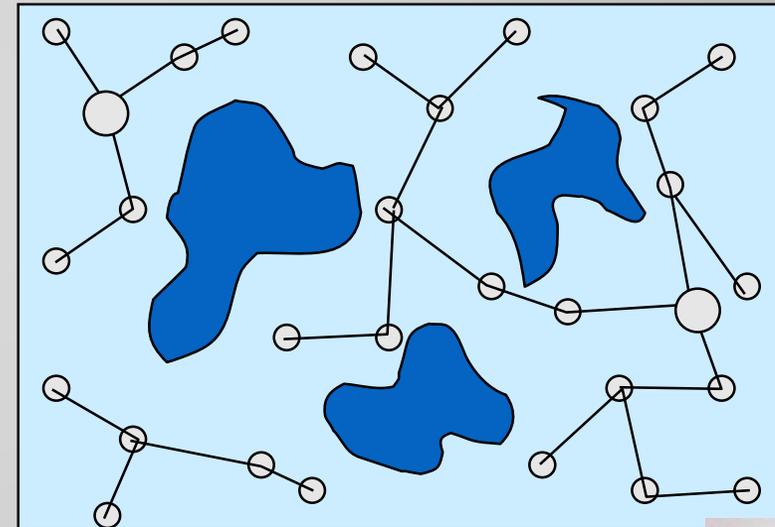
- Collision checker
- Local Method
- Sampler

2 MAJOR STEPS:

- Exploration Phase
- Query Phase

KEY IDEA: EXPLORE RANDOMLY C-SPACE AND CAPTURE C-FREE TOPOLOGY INTO A ROADMAP

COMPLETE IN INFINITE TIME:  
PROBABILISTICALLY COMPLETE



# LOCAL METHOD

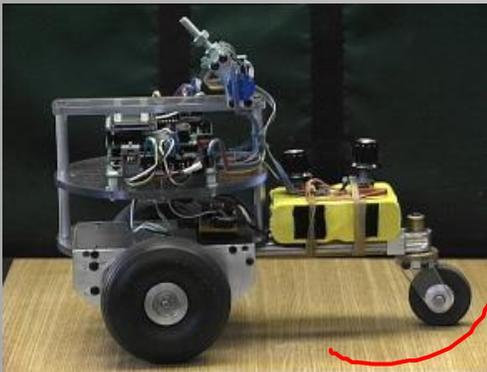
→ IN CHARGE OF KINEMATICS CONSTRAINTS

GIVEN  $q_{INIT}$  AND  $q_{GOAL}$ , HOW TO COMPUTE A LOCAL PATH JOINING THEM?

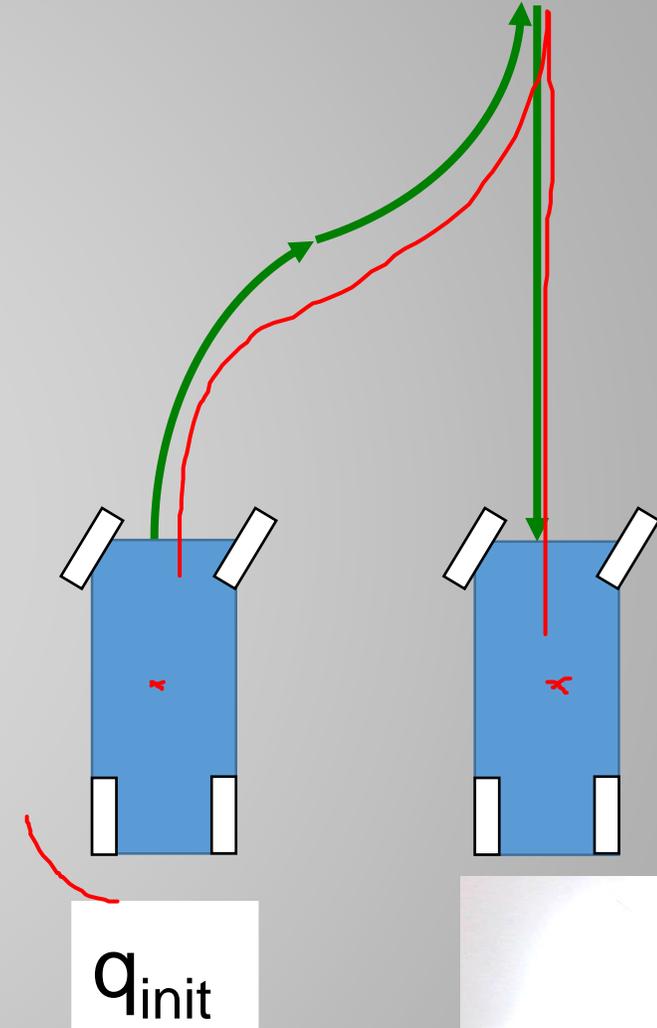
NOT TAKING INTO ACCOUNT OBSTACLES

EXAMPLES:

- Free: linear



Shepp



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

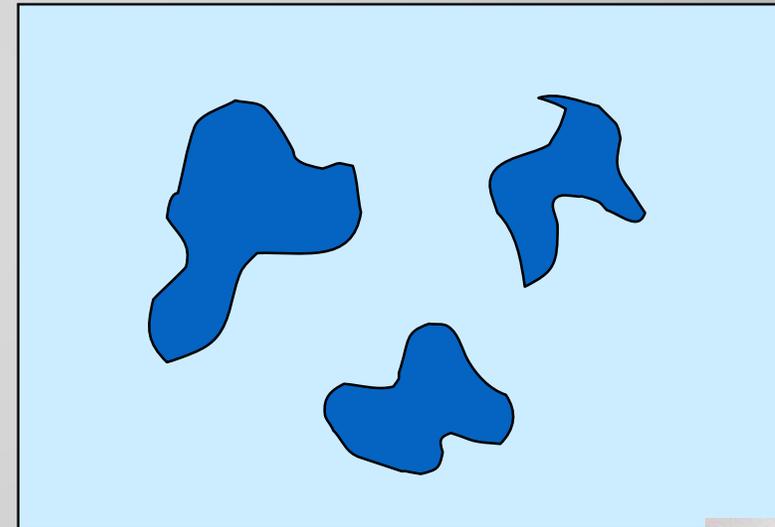
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

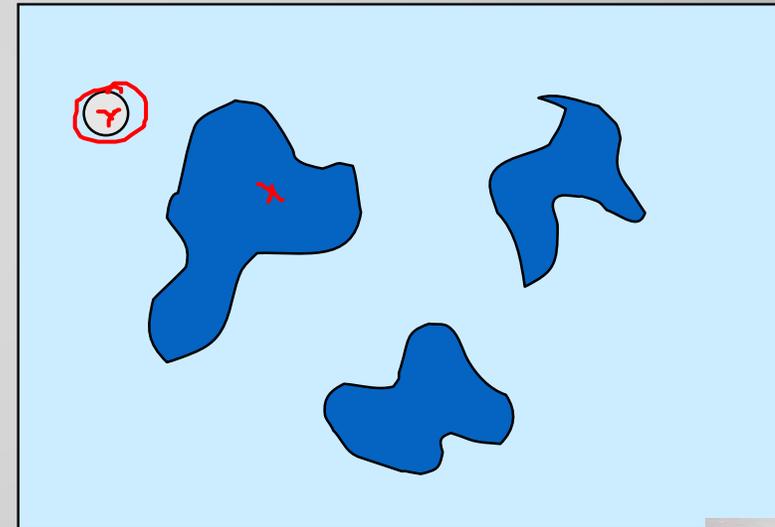
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker ←

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

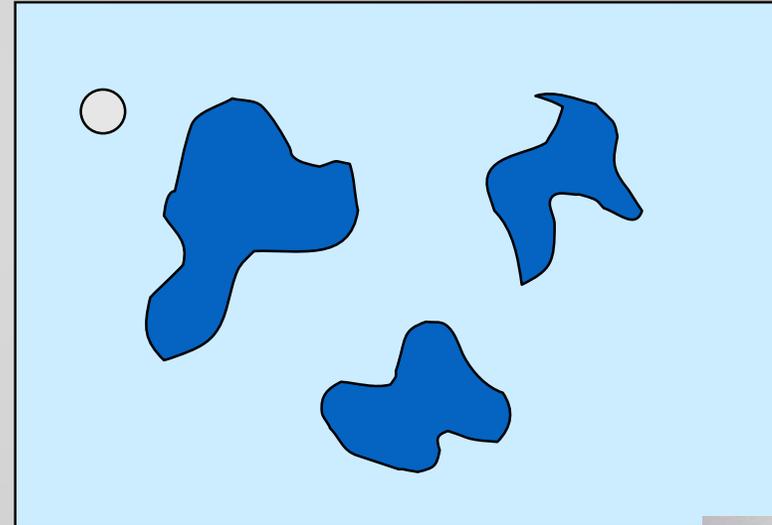
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

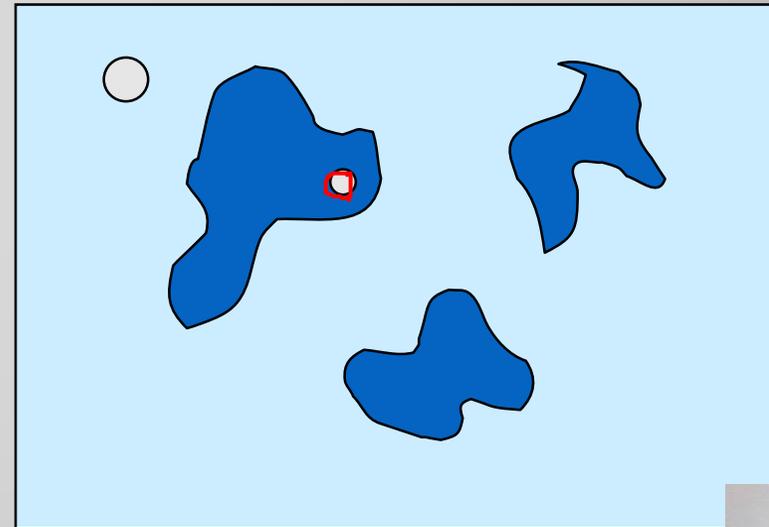
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker 

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

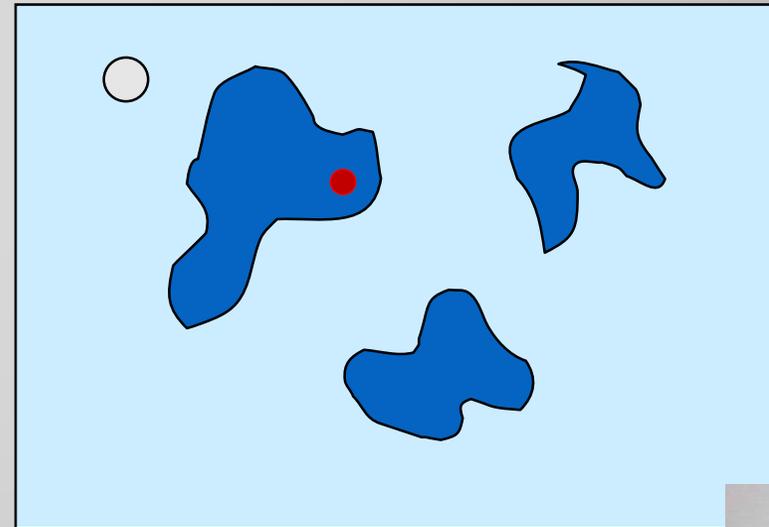
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

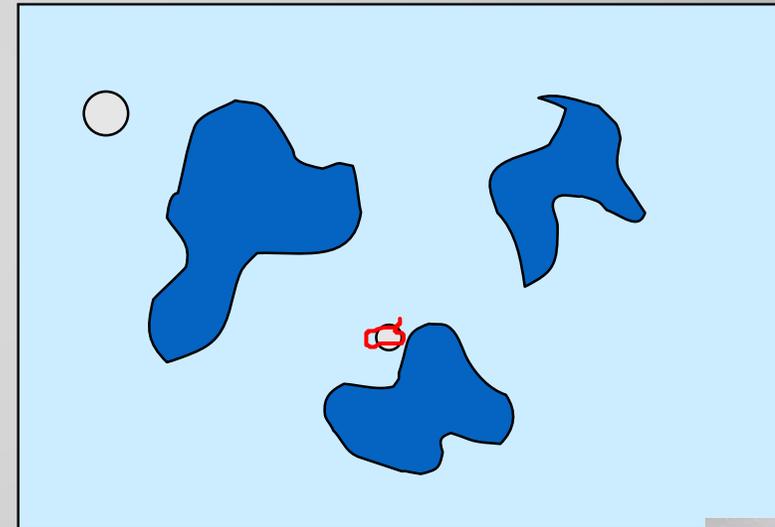
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

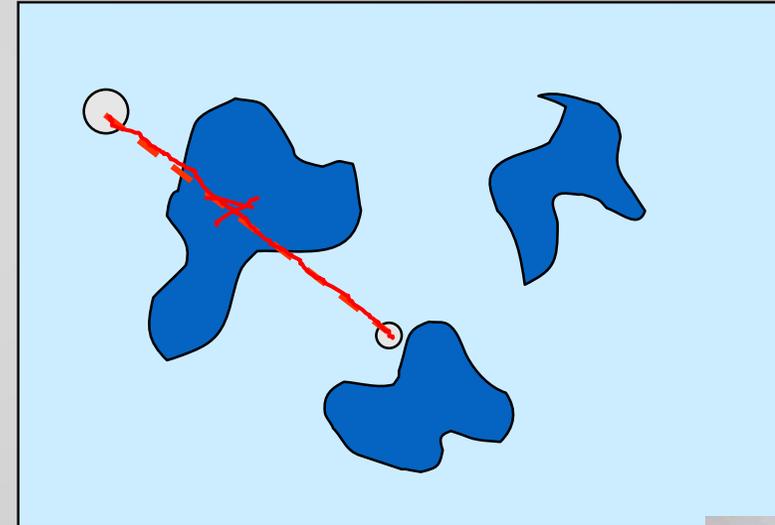
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker ←
- Local method ←←

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

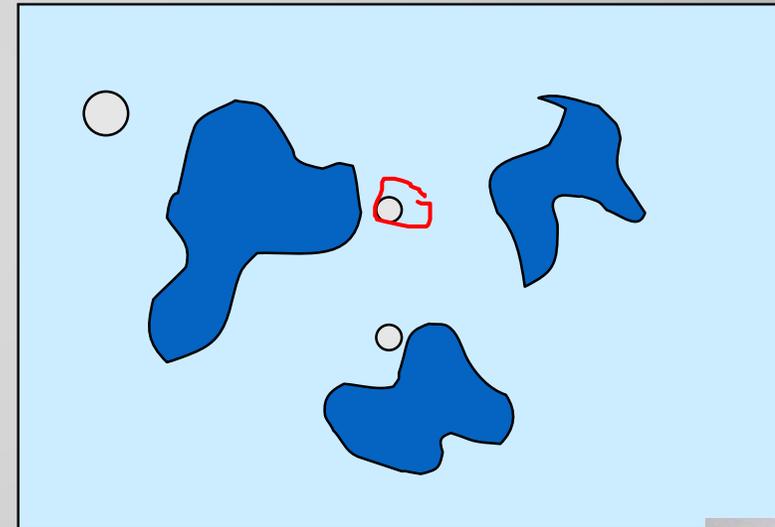
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

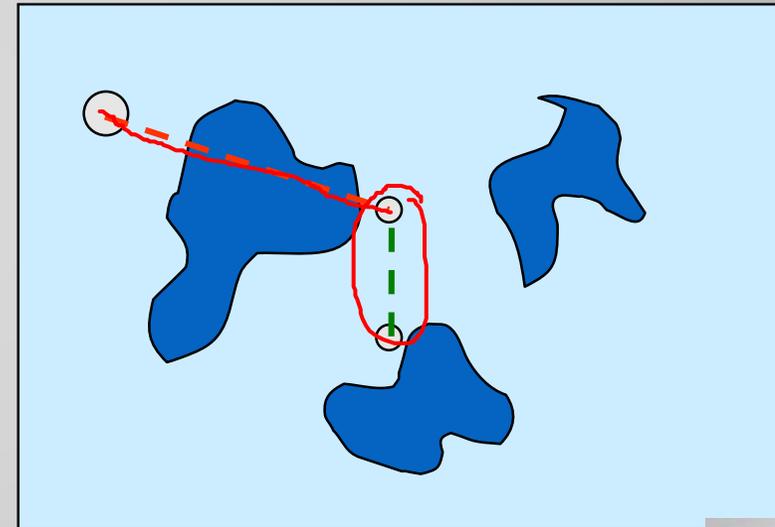
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method —

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

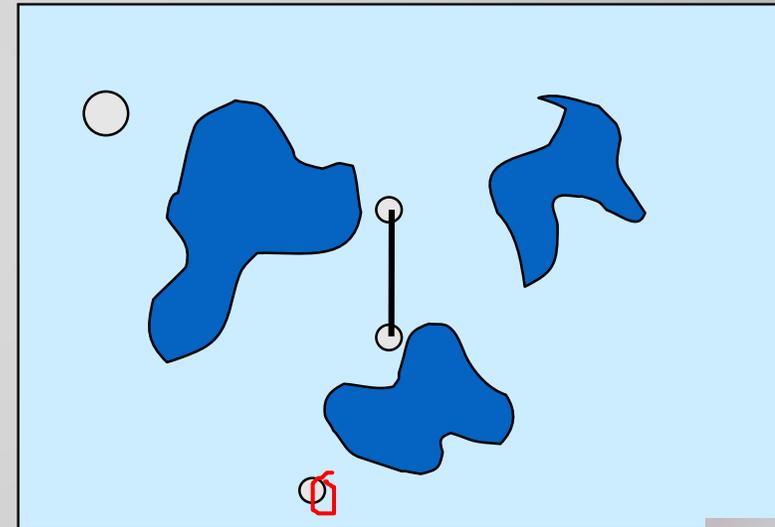
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

### 3. GOTO 1



# PREPROCESSING: LEARNING PHASE

## ITERATIVE ALGORITHM

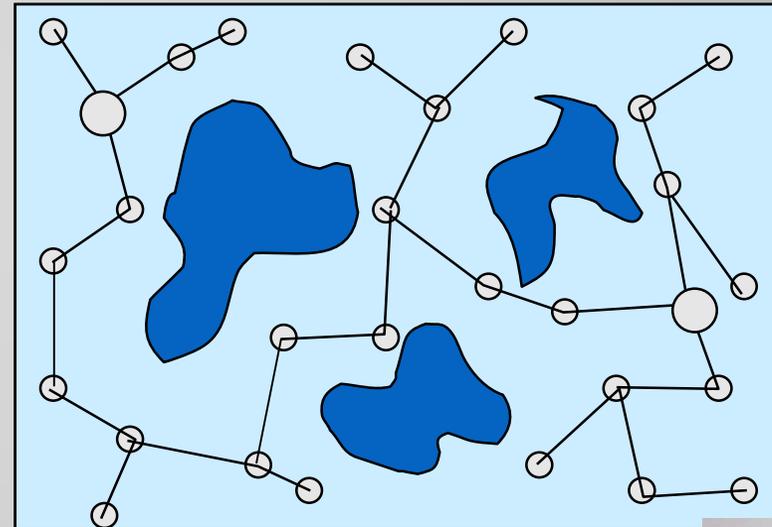
### 1. COMPUTE RANDOM CONFIGURATION

- Collision checker

### 2. CONNECT CONFIGURATION

- Collision checker
- Local method

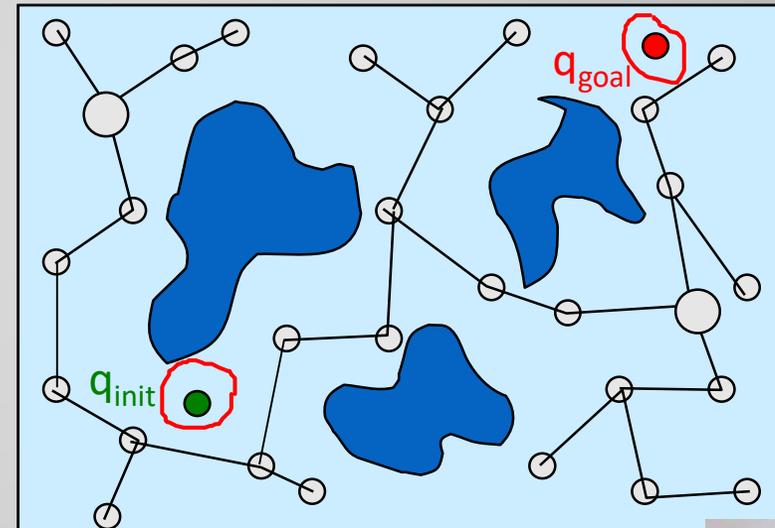
### 3. GOTO 1



# QUERY PHASE

ROADMAP IS REUSED FOR SOLVING QUERIES

1. CONNECT DESIRED INITIAL AND FINAL CONFIGURATIONS
2. IF CORRESPONDING NODES BELONG TO THE SAME CONNECTED COMPONENT, A SOLUTION EXISTS
3. GRAPH SEARCH



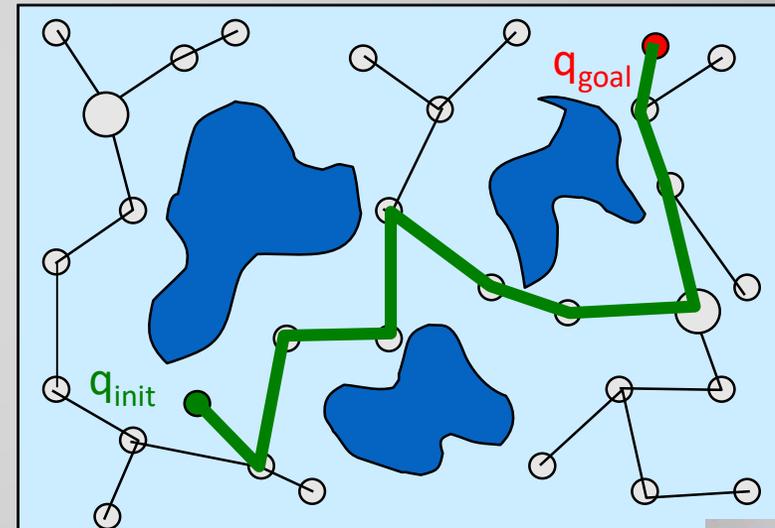
# QUERY PHASE

ROADMAP IS REUSED FOR SOLVING QUERIES

1. CONNECT DESIRED INITIAL AND FINAL CONFIGURATIONS

2. IF CORRESPONDING NODES BELONG TO THE SAME CONNECTED COMPONENT, A SOLUTION EXISTS

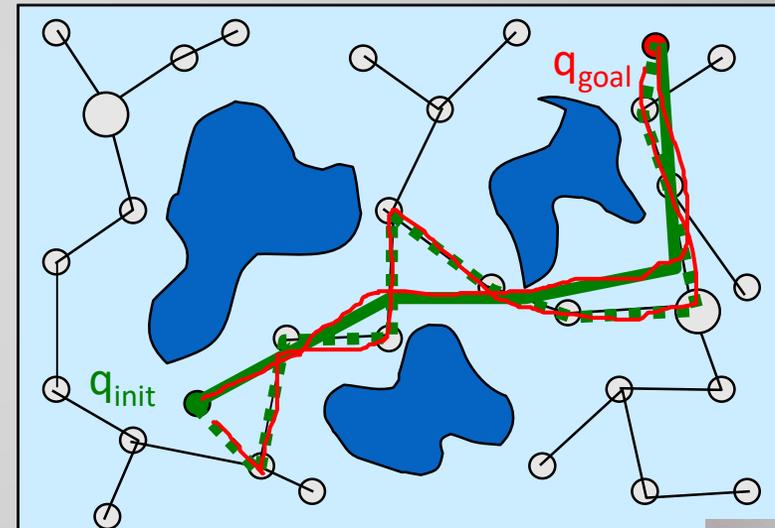
3. GRAPH SEARCH



# QUERY PHASE

ROADMAP IS REUSED FOR SOLVING QUERIES

1. CONNECT DESIRED INITIAL AND FINAL CONFIGURATIONS
2. IF CORRESPONDING NODES BELONG TO THE SAME CONNECTED COMPONENT, A SOLUTION EXISTS
3. GRAPH SEARCH
4. OPTIMIZATION



# GOOD & BAD NEWS

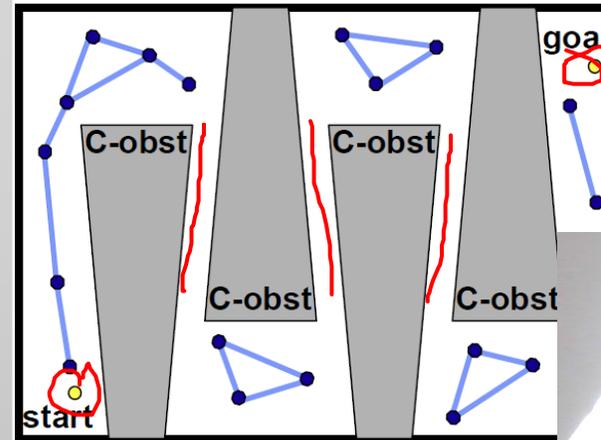
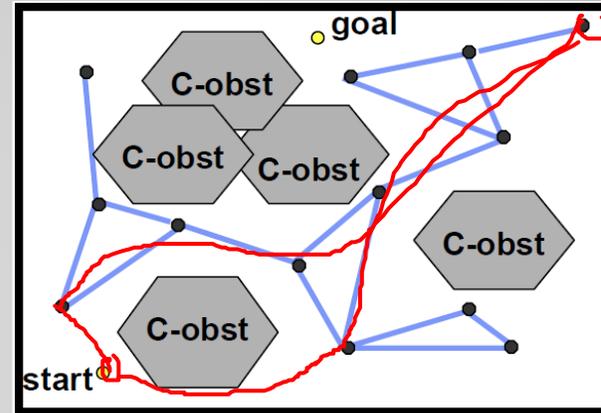
## GOOD NEWS

- Probabilistically complete
- Do not construct the C-space
- Apply easily to high-dimensional C-space
- support fast queries w/ enough preprocessing

MANY SUCCESS STORIES WHERE PRMs SOLVE PREVIOUSLY  
UNSOLVED PROBLEMS

## THE BAD NEWS

- Don't work as well for some problems:
  - ✗ unlikely to sample nodes in *narrow passages*
  - ✗ hard to sample/connect nodes on constraint surfaces
- No optimality or completeness



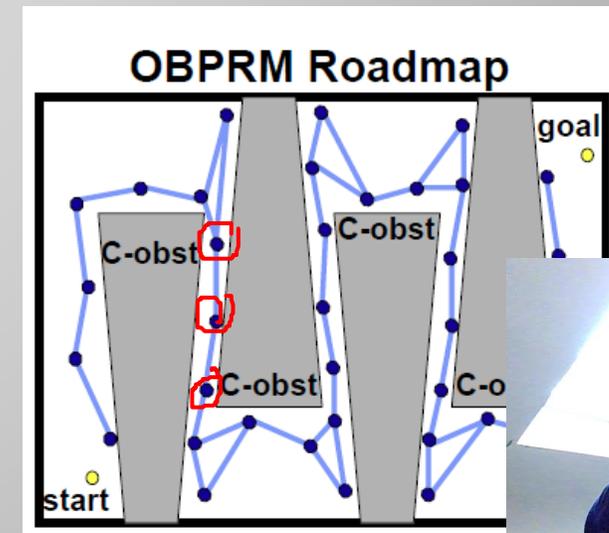
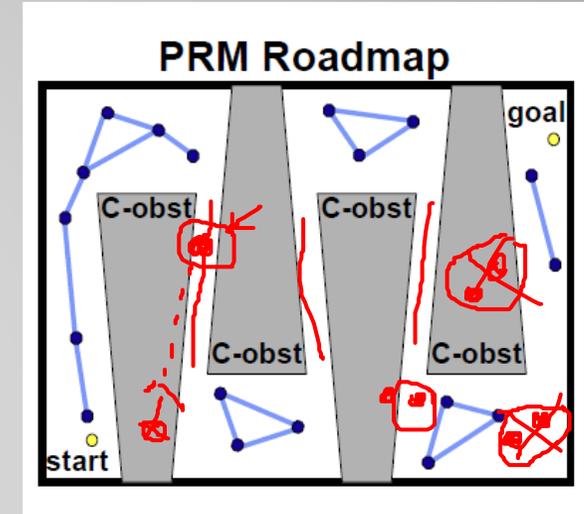
# SAMPLE NEAR OBSTACLES

→ OBPRM

- $q_{in}$  found in collision
- Generate random direction  $v$
- Find  $q_{out}$  in direction  $v$  that is free
- Binary search from  $q_{in}$  to obstacle boundary to generate node

→ GAUSSIAN SAMPLER

- Find a  $q_1$
- Find another  $q_2$  picked from a Gaussian distribution centered at  $q_1$
- If they are both in collision or free, discard. Otherwise, keep the free



# OBPRM: FINDING POINTS ON C-OBSTACLES

## BASIC IDEA:

1. FIND A POINT IN  $S$ 'S C-OBSTACLE (ROBOT PLACEMENT COLLIDING WITH  $S$ )
2. SELECT A RANDOM DIRECTION IN  $C$ -SPACE
3. FIND A FREE POINT IN THAT DIRECTION
4. FIND BOUNDARY POINT BETWEEN THEM USING BINARY SEARCH (COLLISION CHECKS)

**PRM**

- 328 nodes
- 4 major CCs

**OBPRM**

- 161 nodes
- 2 major CCs

RI 16-735, Howie Choset with slides from Nancy Amato, Sujay Bhattacharjee, G.D. Hager, S. LaValle, and a lot from



# SAMPLING STRATEGY

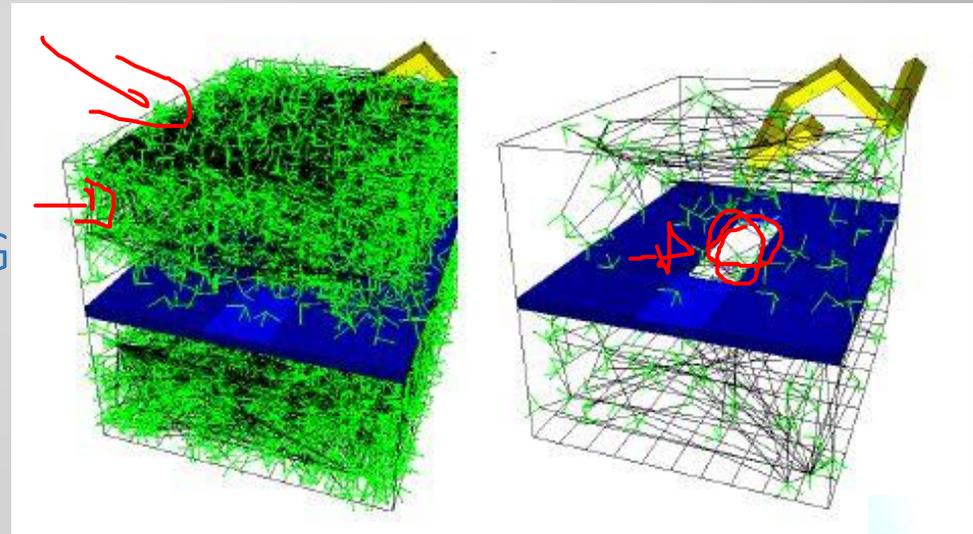
HIGHLY CONSTRAINED PROBLEMS RESULT IN  
HUGE ROADMAPS:

- Construction is time consuming
- Search is time consuming

SAMPLING STRATEGIES HELP IN REDUCING  
THE ROADMAP SIZE

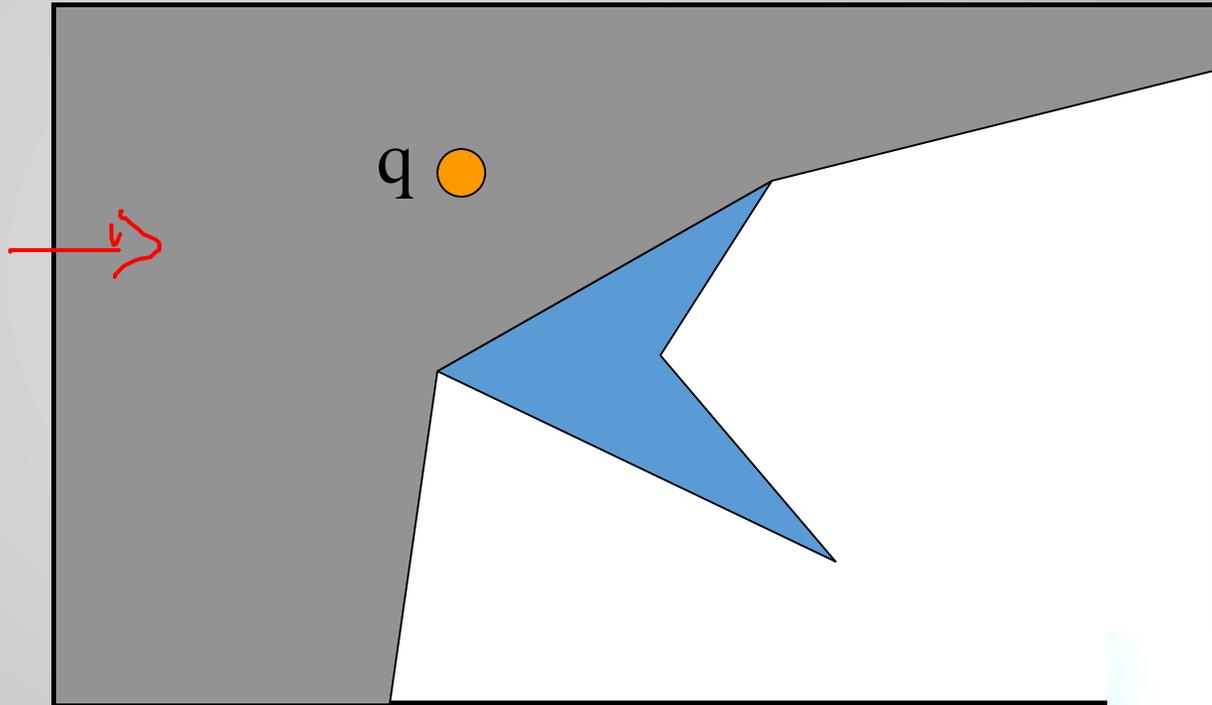
EXAMPLE:

- Visibility-PRM



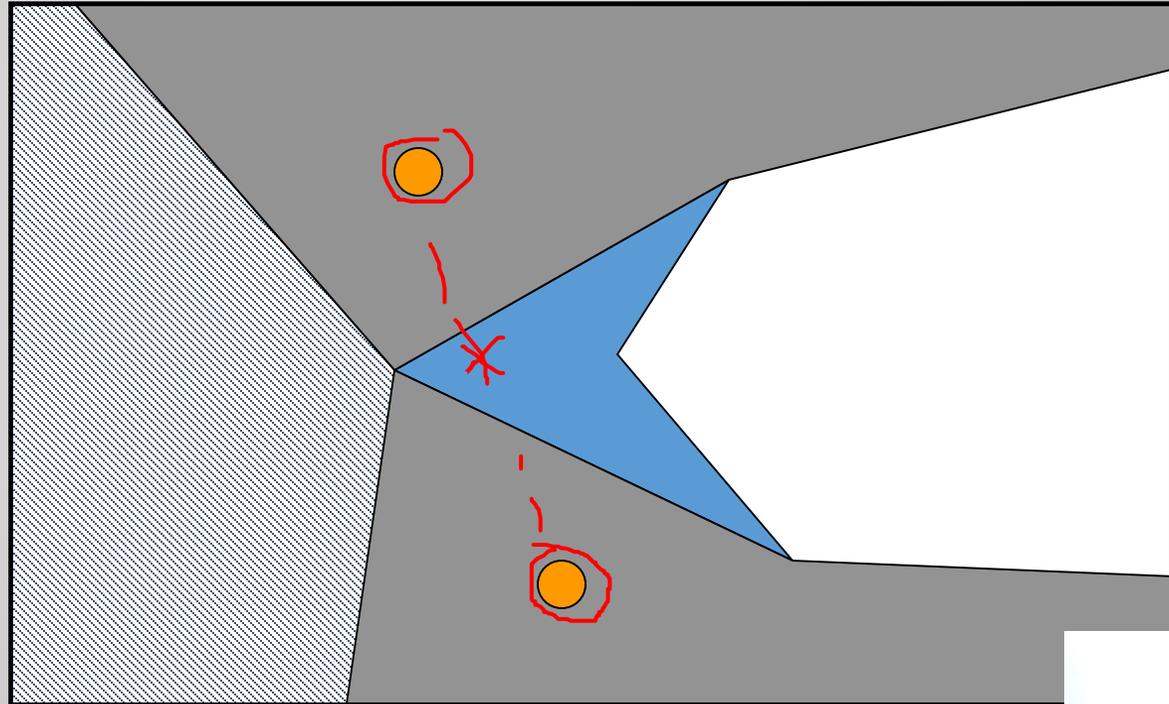
# VISIBILITY-PRM

VISIBILITY DOMAIN OF  
CONFIGURATION Q:



# VISIBILITY-PRM

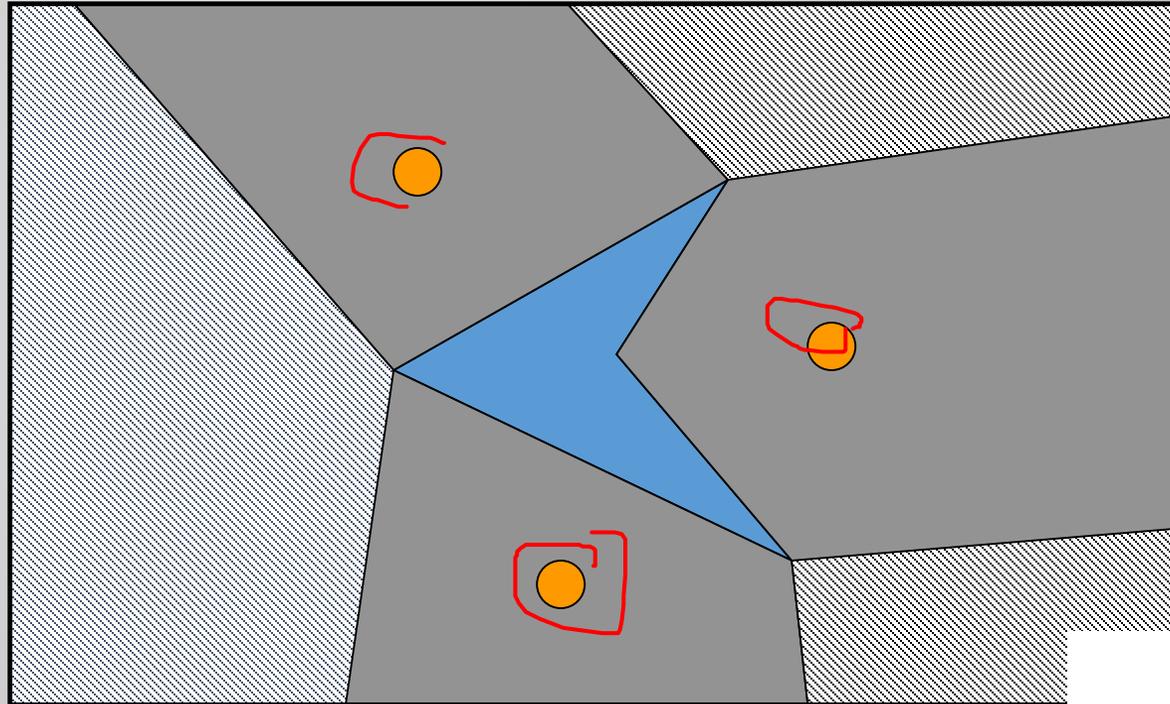
A NEW CONFIGURATION  
IS RETAINED ONLY IF  
OUT OF THE VISIBILITY  
DOMAIN OF OTHER  
CONFIGURATIONS



# VISIBILITY-PRM

A NEW CONFIGURATION IS RETAINED ONLY IF OUT OF THE VISIBILITY DOMAIN OF OTHER CONFIGURATIONS

THESE CONFIGURATIONS ARE CALLED "GUARDIANS"

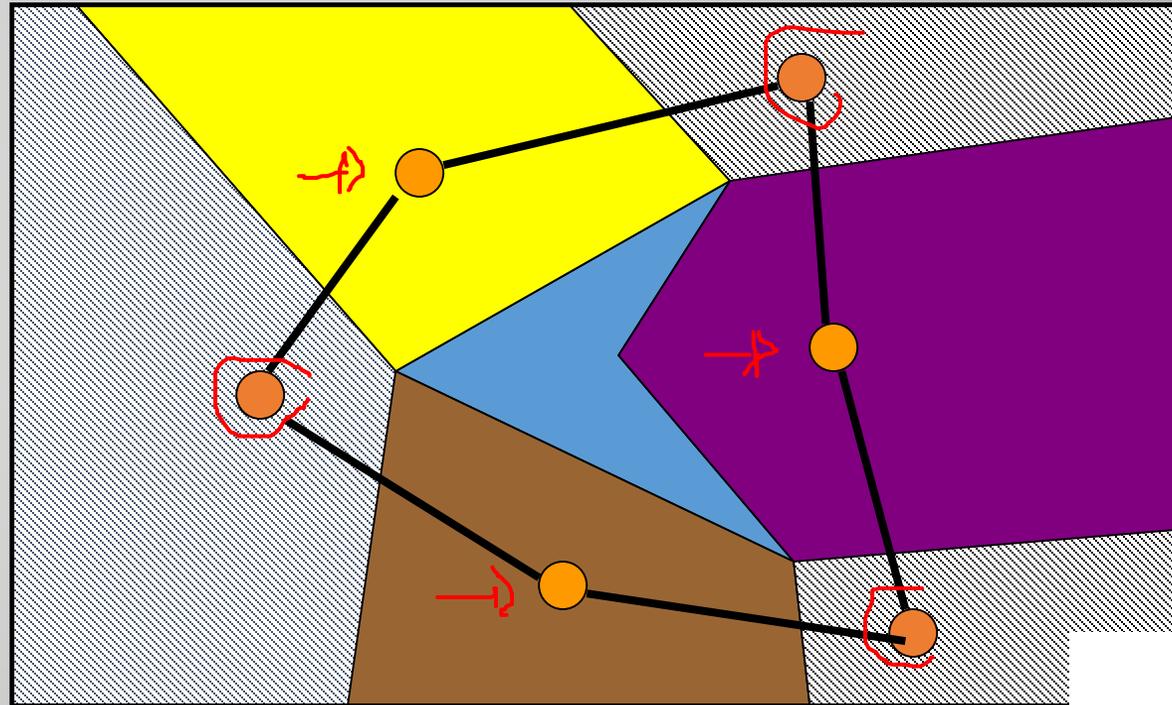


# VISIBILITY-PRM

A NEW CONFIGURATION IS RETAINED ONLY IF OUT OF THE VISIBILITY DOMAIN OF OTHER CONFIGURATIONS

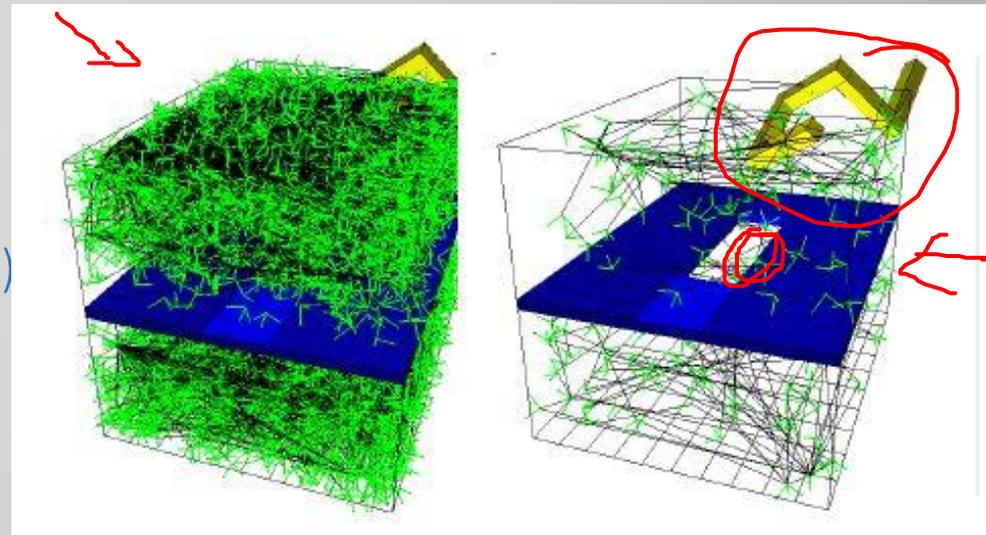
OR IF ALLOW TO CONNECT 2 GUARDIANS

THESE CONFIGURATIONS ARE CALLED "CONNECTORS"



# VISIBILITY-PRM

(THIS IS A 6-DIMENSIONAL C-SPACE IN 3-D)

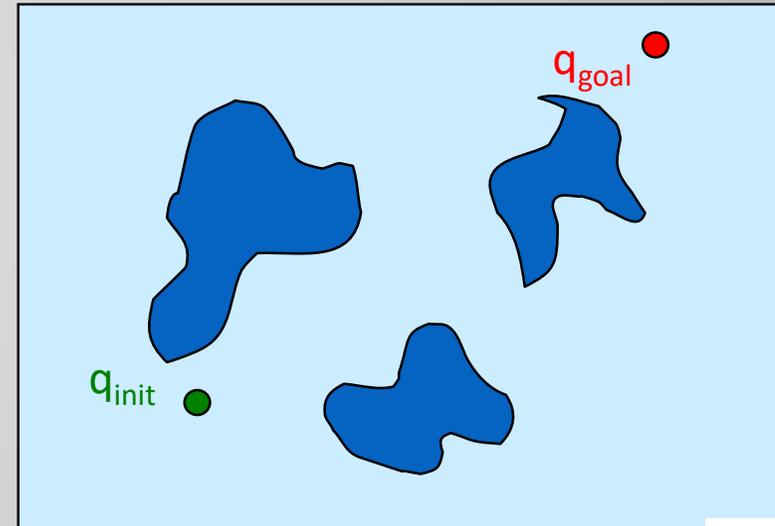




# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

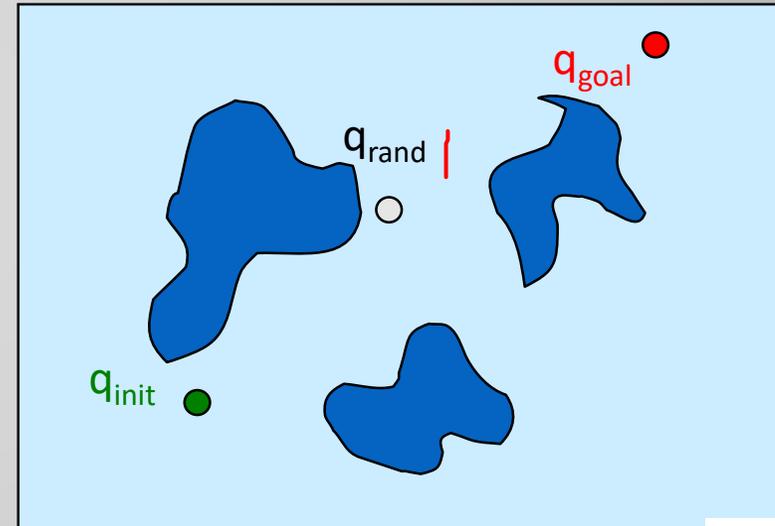
1. COMPUTE  $q_{\text{RAND}}$
2. CONNECT TO  $q_{\text{NEAR}}$
3. INSERT  $q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

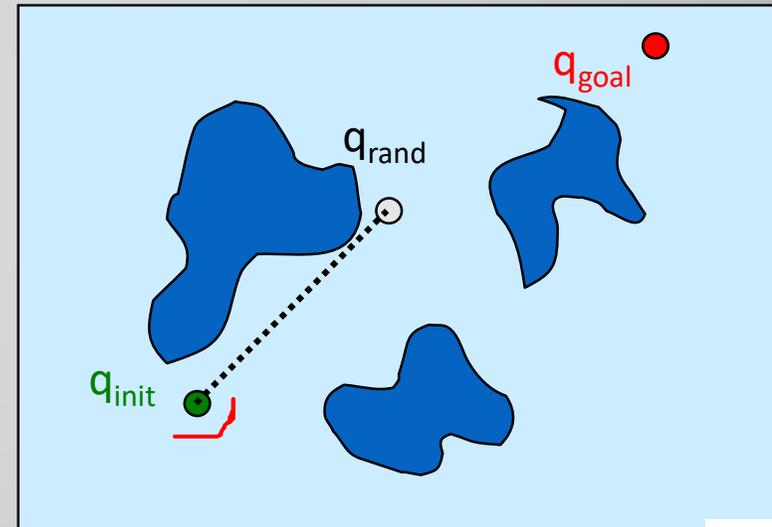
1. COMPUTE  $q_{\text{RAND}}$
2. CONNECT TO  $q_{\text{NEAR}}$
3. INSERT  $q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

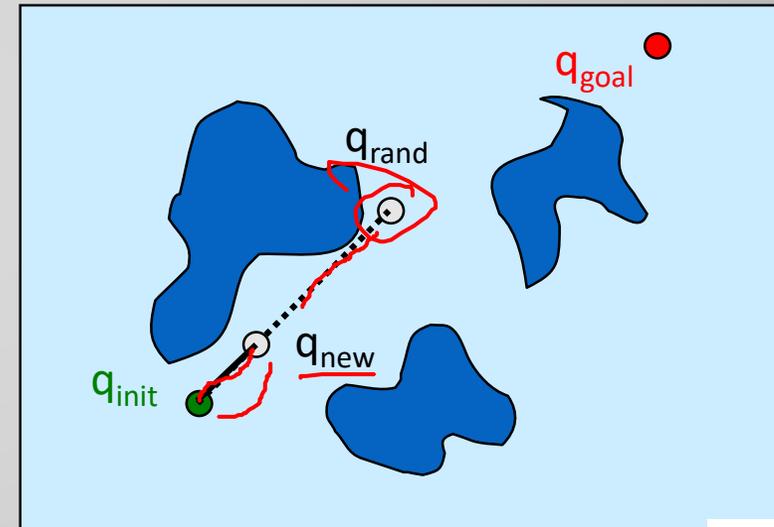
1. COMPUTE  $q_{\text{RAND}}$
2. CONNECT TO  $q_{\text{NEAR}}$
3. INSERT  $q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

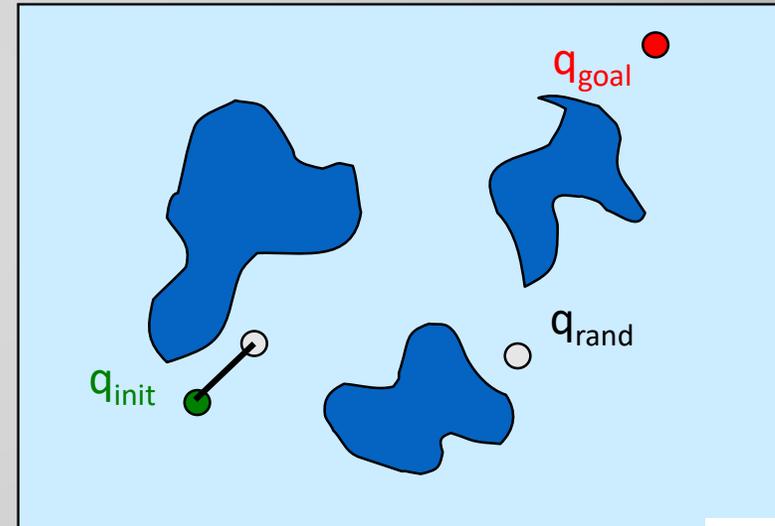
1. COMPUTE  $Q_{\text{RAND}}$
2. CONNECT TO  $Q_{\text{NEAR}}$
3. INSERT  $Q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

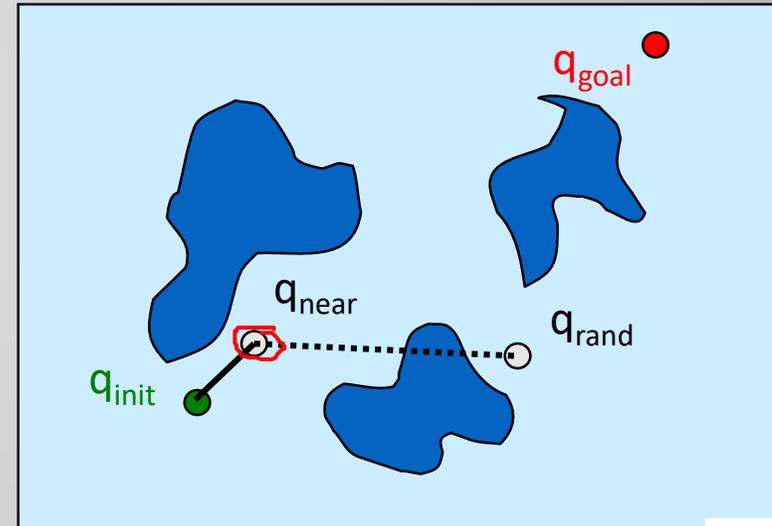
1. COMPUTE  $q_{\text{RAND}}$
2. CONNECT TO  $q_{\text{NEAR}}$
3. INSERT  $q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

ITERATIVE ALGORITHM:

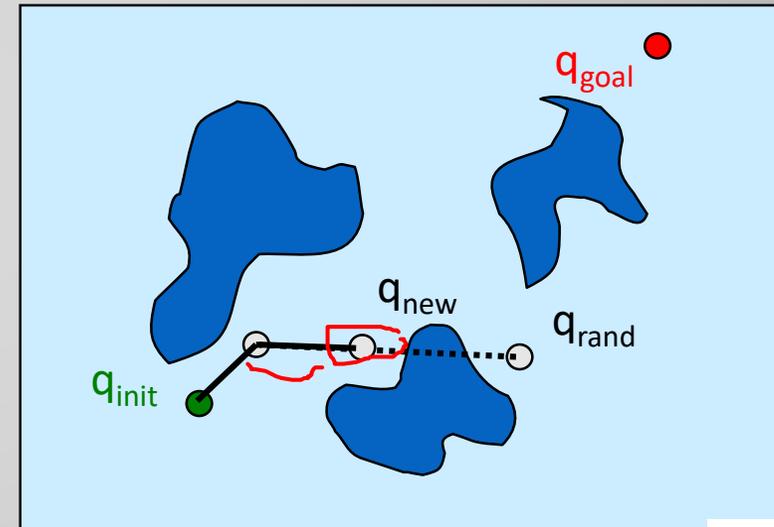
1. COMPUTE  $Q_{\text{RAND}}$
2. CONNECT TO  $Q_{\text{NEAR}}$
3. INSERT  $Q_{\text{NEW}}$
4. GOTO 1



# RRT: RAPIDLY-EXPLORING RANDOM TREES

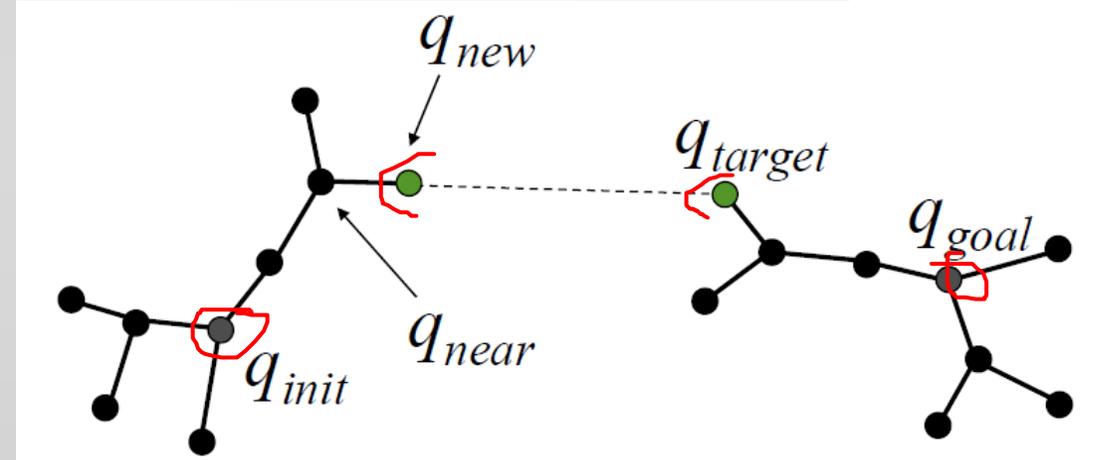
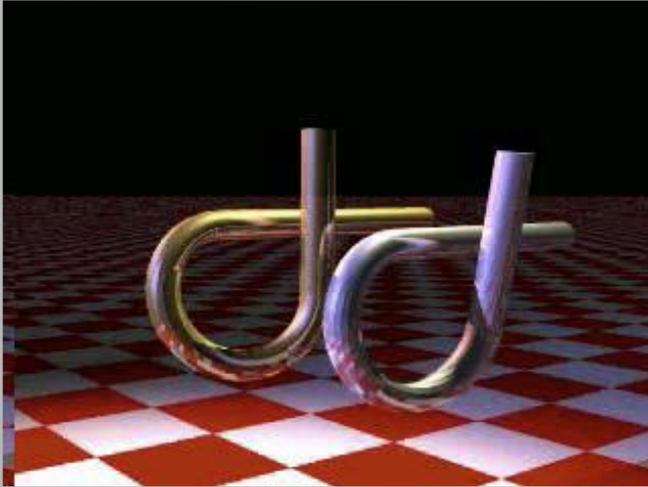
ITERATIVE ALGORITHM:

1. COMPUTE  $Q_{\text{RAND}}$
2. CONNECT TO  $Q_{\text{NEAR}}$
3. INSERT  $Q_{\text{NEW}}$
4. GOTO 1

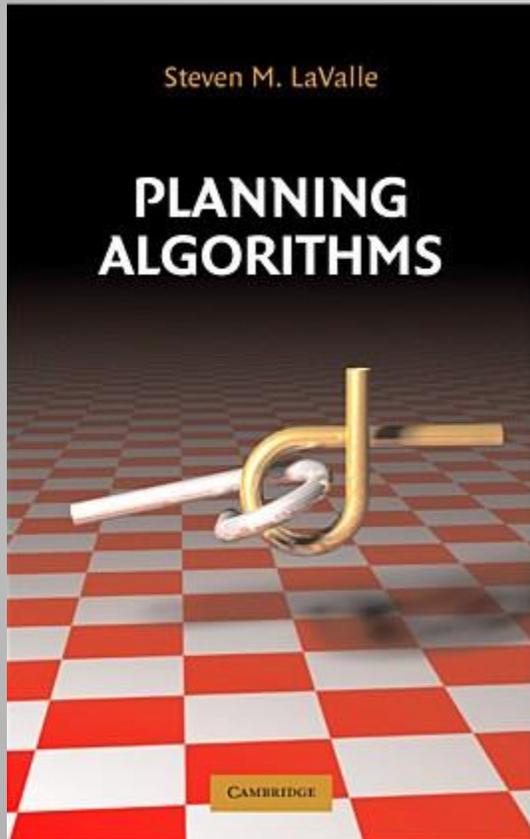




# GROW TWO RRTs TOWARDS EACH OTHER



MORE...



PLANNING ALGORITHMS  
STEVEN M. LAVALLE  
[HTTP://PLANNING.CS.UIUC.EDU/](http://planning.cs.uiuc.edu/)

